# A FAST BOUNDARY ELEMENT METHOD FOR THE SOLUTION OF PERIODIC MANY-INCLUSION PROBLEMS VIA HIERARCHICAL MATRIX TECHNIQUES ${ }^{*, * *, * * *}$ 

Paul Cazeaux ${ }^{1}$ and Olivier Zahm ${ }^{2}$


#### Abstract

Our work is motivated by numerical homogenization of materials such as concrete, modeled as composites structured as randomly distributed inclusions imbedded in a matrix. In this paper, we propose a method for the approximation of the periodic corrector problem based on boundary integral equations. The fully populated matrices obtained by the discretization of the integral operators are successfully dealt with using the $\mathcal{H}$-matrix format.


Résumé. Nous nous intéressons à l'homogénéisation de matériaux de type béton, c'est à dire composés d'agrégats distribués aléatoirement dans une matrice. Nous proposons une méthode pour la résolution du problème de correcteur périodique en s'appuyant sur une formulation intégrale surfacique. Les matrices pleines obtenues par la discrétisation des opérateurs intégraux sont traitées avec succès à l'aide de matrices hiérarchiques.

## Introduction

Understanding the macroscopic properties of composite materials is of great interest in a number of industrial applications. One such example is the study of the aging of electrical nuclear plants due to the long-term behavior of concrete. Homogenization techniques have become widely used to predict this macroscopic behavior, based on the knowledge of the distribution and characteristics of the constituting elements in the microstructure of such materials. A number of different homogenization frameworks exist, ranging from qualitative methods proposed by physicists and engineering, see e.g. [21], to more rigorous mathematical methods developed for example in the periodic case $[6,29]$ or the stationary random case [6,24], giving rise to a vast body of literature.

Using these techniques, the effective (macroscopic) properties of the composite are typically deduced from the solution of an elliptic boundary value problem, called corrector problem. Formulated on a representative volume element (RVE) of the microstructure, solving numerically this problem enables us to sample the local behavior of the microstructure under a macroscopic forcing. Approximate effective parameters are then obtained by averaging the computed local quantities over the RVE [6]. In the case of idealized structures such as perfect crystals, the geometry of the RVE is simple. However, in the case of composite materials such as concrete,

[^0]this geometry can become quite complex. In fact, when the composite is modeled as a matrix containing a stationary random distribution of inclusions, the corrector problem is effectively set on the whole space [24]. In this case, a sequence of corrector problems is usually solved in a bounded box of increasing size. The sequence of approximate effective parameters computed in this way are known to converge to the true effective parameters as the size of the box goes to infinity [14]. Thus it is necessary to develop efficient numerical methods able to deal with complex three-dimensional geometries, induced by large numbers of inclusions in the RVE, see e.g. [9].

A number of numerical methods can be used to approach the solution of the corrector problem. A widely used approach is the direct discretization of the elliptic PDE system, either by finite differences or by the finite elements method. However, such full field simulations remain computationally limited by the large number of degrees of freedom that is required, in particular in a three dimensional setting.

We present a boundary integral reformulation of the corrector problem. We propose a numerical approach for its solution, using the boundary element method (BEM) [23], accelerated by the use of hierarchical matrices. Such methods are known to easily handle jumps in the parameters in a complicated geometry, avoid having to discretize the entire domain in a manner that accurately resolves the geometry of the inclusions, can achieve rapid convergence, and have a robust mathematical foundation [2, 23]. Specific difficulties in this framework involve in particular dealing with:

- periodic boundary conditions,
- a three-dimensional setting,
- complex geometries involving large numbers of randomly distributed inclusions in a representative periodic cell.

While the overwhelming majority of works devoted to the numerical evaluation of effective properties of composite materials employ the finite element method [26], the BEM has been proposed in a few publications as a promising alternative to study the homogenization of composites, see e.g. [12, 16, 25]. A related investigation on periodic problems for the Stokes equation can be found in [17]. In particular, fast multipole methods have recently been designed to increase the efficiency of this approach while dealing with the periodic boundary conditions in the case of rigid inclusions for linear elasticity with applications to the homogenization of carbon nanotube materials in two or three dimensions [22,27], based on the original fast multipole algorithms [18], see also [28].

A closely related problem is the calculation of band structure in periodic crystals or materials, to which boundary integral methods have been recently applied, see e.g. [32]. In particular, our approach is inspired by ideas proposed in [2] in a two-dimensional setting. There is also an extensive literature on the subject of integral equations for scattering from periodic structures, see e.g. [11].

The BEM discretization results in linear systems involving fully populated matrices, by contrast with the sparse systems obtained with the finite differences or finite elements methods. An algebraic tool designed to deal efficiently with such fully populated matrices is the hierarchical matrix (or $\mathcal{H}$-matrix) format introduced by Hackbusch and his collaborators, see e.g. $[5,7,19]$. The main advantages of the $\mathcal{H}$-matrix format are:

- the controlled approximation of the matrix with respect to a given precision,
- the reduced computational and memory cost compared to the usual matrix storage,
- and the accelerated algebraic operations (matrix-vector product, linear solvers, etc).

The structure of the paper is as follows. In section 1.1, we present the boundary integral formulation of the periodic corrector problem. In section 2 , we recall the framework of the $\mathcal{H}$-matrix format by presenting some definitions, the main properties and associated algorithms. Finally, in section 3 we present some numerical results on the application of the $\mathcal{H}$-matrix technique to the corrector problem.

## 1. The periodic corrector problem

### 1.1. Reformulation of the corrector problem as a boundary integral problem

In this section we derive a boundary integral formulation for the corrector problem associated to the Laplace equation. Let $\Omega=(-1 / 2,1 / 2)^{d}$ be an open unit periodic cell in $\mathbb{R}^{d}, d=2,3$. We decompose $\Omega$ as a set of closed inclusions $\Omega_{\text {int }} \subset \Omega$ and a connected matrix $\Omega_{\text {ext }}=\Omega \backslash \overline{\Omega_{\text {int }}}$, see figure 1. In the following, $\Gamma=\partial \Omega_{\text {int }}$ denotes the boundary of the set of inclusions.


Figure 1. Example of the geometry of a set of inclusions.

Remark 1.1. It is not necessary to suppose that the inclusions are distributed strictly inside the periodic cell.
Let $\kappa(\mathbf{x})$ be a scalar diffusion coefficient which takes two different values $\kappa_{\text {int }}, \kappa_{\text {ext }}>0$ respectively in the inclusions and in the matrix:

$$
\kappa(\mathbf{x})= \begin{cases}\kappa_{i n t} & \text { in } \Omega_{i n t}  \tag{1}\\ \kappa_{e x t} & \text { in } \Omega_{e x t}\end{cases}
$$

We are interested in the periodic corrector $u \in H^{1}(\Omega) / \mathbb{R}$ solution of the following boundary value problem:

$$
\left\{\begin{array}{c}
-\operatorname{div}(\kappa(\mathbf{E}+\nabla u))=0, \quad \text { in } \Omega,  \tag{2}\\
u \text { is } \Omega \text {-periodic },
\end{array}\right.
$$

where $\mathbf{E} \in \mathbb{R}^{3}$ is a given macroscopic gradient of some diffusive quantity $\widetilde{u}$. The problem (2) is the system of equilibrium equations on the RVE for the Laplace equation, assuming that the real gradient is split into two parts as

$$
\nabla \widetilde{u}=\mathbf{E}+\nabla u
$$

where $\mathbf{E}$ is the overall gradient and $\nabla u$ the fluctuating terms. It is well-known that problem (2) has a unique solution, see e.g. [1]. Note that the solution $u$ is harmonic in $\Omega \backslash \Gamma$ and satisfies continuity conditions:

$$
\begin{equation*}
u, \kappa(\nabla u+\mathbf{E}) \cdot \mathbf{n}^{e x t} \text { continuous across } \Gamma \tag{3}
\end{equation*}
$$

where $\mathbf{n}^{e x t}$ is the exterior normal to $\Gamma$. An elegant approach for the integral representation of periodic fields involves the use of the Green's function satisfying the periodic boundary conditions exactly. Let $\delta_{0}$ be the Dirac
delta function centered at the origin. The periodic Green's function $G_{p e r}$ is then defined by $-\Delta G_{p e r}=\delta_{0}-1$ over $\Omega$ (note that the right-hand side must have zero mean over the periodic cell). Let us introduce the single-layer potential [23] as an operator $H^{-1 / 2}(\Gamma) \rightarrow H^{1}(\Omega)$ defined by:

$$
\begin{equation*}
\left(\Psi_{S L} \sigma\right)(\mathbf{x})=\int_{\Gamma} G_{p e r}(\mathbf{x}-\mathbf{y}) \sigma(\mathbf{y}) \mathrm{d} s(\mathbf{y}) \quad \text { for } \mathbf{x} \in \Omega \backslash \Gamma \tag{4}
\end{equation*}
$$

We can represent the solution to the PDE problem (2) as $u=\Psi_{S L} \sigma$ with a single-layer density $\sigma$. It remains to solve for the density $\sigma \in H^{-1 / 2}(\Gamma)$ so that the matching conditions (3) are satisfied. We use the standard jump relations for single layer potential [23] to write

$$
\begin{equation*}
\kappa_{e x t}\left(-1 / 2 I d+K_{0}^{\prime}\right) \sigma-\kappa_{i n t}\left(1 / 2 I d+K_{0}^{\prime}\right) \sigma=-\left(\kappa_{e x t}-\kappa_{i n t}\right) \mathbf{E} \cdot \mathbf{n}^{e x t} \quad \text { on } \Gamma, \tag{5}
\end{equation*}
$$

where $I d$ is the identity operator and $K_{0}^{\prime}$ is the adjoint double layer boundary integral operator $H^{-1 / 2}(\Gamma) \rightarrow$ $H^{1 / 2}(\Gamma)$ formally defined as:

$$
\begin{equation*}
\left(K_{0}^{\prime} \sigma\right)(\mathbf{x})=\int_{\Gamma} \sigma(\mathbf{y}) \nabla_{\mathbf{x}} G_{p e r}(\mathbf{x}-\mathbf{y}) \cdot \mathbf{n}^{e x t}(\mathbf{y}) \mathrm{d} s(\mathbf{y}) \tag{6}
\end{equation*}
$$

Finally, we obtain the following boundary integral equation:

$$
\begin{equation*}
\left(\frac{\kappa_{e x t}+\kappa_{i n t}}{\kappa_{e x t}-\kappa_{i n t}} I d-K_{0}^{\prime}\right) \sigma=\mathbf{E} \cdot \mathbf{n}^{e x t} \tag{7}
\end{equation*}
$$

where the operator appearing on the left-hand side is a compact perturbation of the identity.

### 1.2. Approximation of the periodic Green's function

A Galerkin approach is used to discretize (7) given a mesh of the boundary $\Gamma$. Functions in $H^{-1 / 2}(\Gamma)$ are approximated by constant-by-cell functions. The main difficulty is the computation of values of the periodic Green's function necessary to the discretization of the boundary integral operator $K_{0}^{\prime}$ defined by (6). Indeed, while the free-space Green's function is given by an analytical expression, the periodic Green's function is not known explicitely. It is rather given by a Fourier series,

$$
G_{p e r}(\mathbf{x})=\sum_{\mathbf{k} \in \mathbb{Z}^{d} ; \mathbf{k} \neq 0} \frac{e^{2 i \pi \mathbf{k} \cdot \mathbf{x}}}{|\mathbf{k}|^{2}}
$$

Because such sums converge too slowly to be numerically useful, a number of schemes have been devised for their evaluation, see e.g. [8,13]. We use here a different, new approach, based on an idea proposed recently by Barnett and Greengard [2] for the Helmholtz two-dimensional case. Observe that the periodic Green's function can be represented in $\Omega$ as

$$
\begin{equation*}
G_{p e r}(\mathbf{x})=G_{\infty}(\mathbf{x})+\frac{|\mathbf{x}|^{2}}{6}+G_{p e r}^{r}(\mathbf{x}) \tag{8}
\end{equation*}
$$

where $G_{\infty}$ is the free-space Green's function (satisfying $-\Delta G_{\infty}=\delta_{0}$ over the whole domain $\mathbb{R}^{d}$ ), and $G_{p e r}^{r}$, the "regular" part of the periodic Green's function, is a smooth, harmonic function in $\Omega$. It can thus be expanded on the basis of the solid spherical harmonics as a series converging uniformly in $\Omega$ :

$$
\begin{equation*}
G_{p e r}^{r}(\mathbf{x})=\sum_{l=0}^{\infty} \sum_{m=-l}^{l} \alpha_{l}^{m} \Phi_{l}^{m}(\mathbf{x}) \tag{9}
\end{equation*}
$$

where $\Phi_{l}^{m}$ is the regular solid harmonics of degree $l$ and order $m$ and the $\alpha_{l}^{m}$ are scalar coefficients. An approximation to $G_{p e r}^{r}$ can then be computed by truncating the expression (9) up to some degree $L \in \mathbb{N}$, given that we tabulate beforehand the coefficients of the expansion $\alpha_{l}^{m}$ for $l \leq L$.


Figure 2. Convergence in pointwise error, sampled in $\Omega$, for the periodic Green's function scheme.

To achieve this, we construct a small linear system that enables us to compute an approximation to the coefficients $\alpha_{l}^{m}$ easily, as in [2], paragraph 3.2. The general idea is to enforce numerically the periodic boundary conditions on the approximation to $G_{p e r}$ based on the representation (8) and (9). In practice, we use a leastsquares algorithm, minimizing the $L^{2}$-norm of the deviation from periodicity of the function and its normal derivative on $\partial \Omega$, sampled at Gaussian quadrature points. We refer to the forthcoming paper [10] for a more complete description and analysis of this method.

This method, implemented in Matlab, allows us to achieve an exponential convergence rate as shown by the green curve in figure 2. It can be further accelerated by removing from the regular part $G_{p e r}^{r}$ copies of the free-space Green's function centered in the nearest neighbor cells. We then compute and use the coefficients $\beta_{l}^{m}$ in the following expansion:

$$
\begin{equation*}
\widetilde{G_{p e r}^{r}}(\mathbf{x})=\sum_{l=0}^{\infty} \sum_{m=-l}^{l} \beta_{l}^{m} \Phi_{l}^{m}(\mathbf{x})=G_{p e r}(\mathbf{x})-\sum_{\mathbf{m} \in\{-1,0,1\}^{3}} G_{\infty}\left(\mathbf{x}+\sum_{i=1}^{3} m_{i} \mathbf{e}_{i}\right)-\frac{|\mathbf{x}|^{2}}{6} \tag{10}
\end{equation*}
$$

where the $\mathbf{e}_{i}$ are the vectors of the standard basis of $\mathbb{R}^{3}$. The convergence rate is much improved as seen in figure 2. Note that, due to the symmetry of the square unit cell, only harmonics of even degree $l \geq 4$ and order $m$ multiple of 4 have a nonzero coefficient $\beta_{l}^{m}$ in (10). We use this approximation below, fixing $L=8$, which ensures a relative error in the computation of the corrective term $\widetilde{G_{p e r}^{r}}$ of about $10^{-4}$.

To solve numerically the boundary integral equation (7), we propose to use the classical Boundary Element Method. Note that this approach gives rise to full matrices for the representation of the operators, and to
reduce the complexity of the computation, compression techniques such as multipole expansions (e.g. the fast multipole method) or hierarchical matrices can be used; in this paper, this latter approach is investigated (see section 2).

Remark 1.2. Since the goal of this study is to explore the use of the hierarchical matrix format in the context of the solution of the corrector problem, we do not fully detail here the numerical discretization method which is quite standard, see e.g. [30] and we refer to a forthcoming paper [10] for the full description of the approach. We limit ourselves to the case of a smooth boundary $\Gamma$ in dimension $d=3$.

## 2. $\mathcal{H}$-matrix

In this section we recall the main properties of the hierarchical matrix format, also called $\mathcal{H}$-matrix. This format, first introduced by Hackbusch in [19], is an algebraic tool designed mainly to manage large linear systems with fully populated matrices. For a more in-depth description, we refer to [5, 7]. This format as two main advantages for numerical computation:

- it provides a controlled approximation of the matrix that reduces the memory requirements,
- all algebraic operations (matrix-vector product, LU-factorization, etc) can be accelerated compared to the full storage.


### 2.1. Motivations

The $\mathcal{H}$-matrix format can be introduced as a data-sparse approximation of matrices resulting from the discretization of non local integral operators [20] of type:

$$
\begin{equation*}
\int_{\Omega} g(\mathbf{x}, \mathbf{y}) u(\mathbf{y}) \mathrm{d} \mathbf{y}=f(\mathbf{x}), \quad \forall \mathbf{x} \in \Omega \tag{11}
\end{equation*}
$$

where the kernel $g$ (possibly singular) is assumed to be asymptotically smooth, that is to satisfy :

$$
\begin{equation*}
\left|\partial_{\mathbf{x}}^{\alpha} \partial_{\mathbf{y}}^{\beta} g(\mathbf{x}, \mathbf{y})\right| \leq C_{1}\left(C_{2}\|\mathbf{x}-\mathbf{y}\|\right)^{-|\alpha|-|\beta|}|g(\mathbf{x}, \mathbf{y})|, \quad C_{1}, C_{2} \in \mathbb{R} \quad \text { with } \alpha, \beta \in \mathbb{N}^{d} \tag{12}
\end{equation*}
$$

In the present paper, the considered integral equation is the boundary integral equation (7). We look for the weak solution $u \in V$ (with $V$ an appropriate Hilbert space) of problem (11) which satisfies $a(u, v)=b(v)$ for all $v \in V$, where $a(\cdot, \cdot)$ and $b(\cdot)$ are defined by :

$$
\begin{aligned}
a(u, v) & =\int_{\Omega} \int_{\Omega} g(\mathbf{x}, \mathbf{y}) v(\mathbf{x}) u(\mathbf{y}) \mathrm{d} \mathbf{y} \mathrm{~d} \mathbf{x} \\
b(v) & =\int_{\Omega} f(\mathbf{x}) v(\mathbf{x}) \mathrm{d} \mathbf{x}
\end{aligned}
$$

Under classical assumptions on the kernel $g(\cdot, \cdot)$, the bilinear form $a$ is coercive and (by the Lax-Milgram lemma) the variational problem is well posed [20]. The Galerkin approximation $u^{h}$ on a finite element subspace $V^{h}=\operatorname{span}\left\{\phi_{i}\right\}_{1 \leq i \leq n}$ ( $\phi_{i}$ being shape functions with compact, localized support) is defined by the relation $u^{h}=\sum_{i=1}^{n} U_{i} \phi_{i}$, where the coefficients $\left\{U_{i}\right\}_{1 \leq i \leq n}$ are the solution of the linear system

$$
\begin{equation*}
A U=B \tag{13}
\end{equation*}
$$

with $A_{i, j}=a\left(\phi_{i}, \phi_{j}\right)$ and $B_{i}=b\left(\phi_{i}\right)$.
In order to motivate the following, let us consider two subsets of indices $\tau, \sigma \subset\{1, \ldots, n\}$, and $\Omega_{\tau}=$ $\cup_{i \in \tau} \operatorname{supp}\left(\phi_{i}\right)$ and $\Omega_{\sigma}=\cup_{i \in \sigma} \operatorname{supp}\left(\phi_{i}\right)$ two clusters of $\Omega$ such that :

$$
\begin{equation*}
\min \left\{\operatorname{diam}\left(\Omega_{\tau}\right), \operatorname{diam}\left(\Omega_{\sigma}\right)\right\} \leq \eta \operatorname{dist}\left(\Omega_{\tau}, \Omega_{\sigma}\right) \tag{14}
\end{equation*}
$$

with $\eta>0$. This last condition is called the admissibility condition.


Figure 3. Representation of two admissible clusters $\Omega_{\tau}$ and $\Omega_{\sigma}$ (left) and the corresponding block in the matrix $A$ (right).

In such a situation (see figure 3), the restriction of the kernel $g$ on $\Omega_{\tau} \times \Omega_{\sigma}$ usually has a low rank structure: for any $\varepsilon>0$ there exists $k \in \mathbb{N}$, depending on $\varepsilon$ and relatively small [15], and there exists pairs of functions $\left(h_{l}^{1}, h_{l}^{2}\right) \in L^{2}\left(\Omega_{\tau}\right) \times L^{2}\left(\Omega_{\sigma}\right)$ with $l \in\{1, \ldots, k\}$ such that the $\operatorname{sum} \widetilde{g}(\mathbf{x}, \mathbf{y})=\sum_{l=1}^{k} h_{l}^{1}(\mathbf{x}) h_{l}^{2}(\mathbf{y})$, called a $k$-rank function, is such that:

$$
\begin{equation*}
\|g(\mathbf{x}, \mathbf{y})-\widetilde{g}(\mathbf{x}, \mathbf{y})\|_{L^{2}\left(\Omega_{\tau} \times \Omega_{\sigma}\right)} \leq \varepsilon\|g(\mathbf{x}, \mathbf{y})\|_{L^{2}\left(\Omega_{\tau} \times \Omega_{\sigma}\right)} \tag{15}
\end{equation*}
$$

For example, $\widetilde{g}$ can be expressed e.g. with a truncated Taylor series or an interpolation scheme, see [5,7]. The existence of such a low rank approximation ensures that the block $A_{\tau, \sigma}$ also has a low rank structure. Indeed, replacing $g$ by $\widetilde{g}$ in the definition of the block $A_{\tau, \sigma}$, we obtain:

$$
\begin{aligned}
\left(A_{\tau, \sigma}\right)_{i, j} & \approx \int_{\Omega} \int_{\Omega} \widetilde{g}(\mathbf{x}, \mathbf{y}) \phi_{\sigma(j)}(\mathbf{x}) \phi_{\tau(i)}(\mathbf{y}) \mathrm{d} \mathbf{y d} \mathbf{x} \\
& =\sum_{n=1}^{k} \underbrace{\left(\int_{\Omega} h_{n}^{1}(\mathbf{x}) \phi_{\sigma(j)}(\mathbf{x}) \mathrm{d} \mathbf{x}\right)}_{B_{i, k}} \underbrace{\left(\int_{\Omega} h_{n}^{2}(\mathbf{y}) \phi_{\tau(i)}(\mathbf{y}) \mathrm{d} \mathbf{y}\right)}_{C_{j, k}}=\left(B C^{t}\right)_{i, j}
\end{aligned}
$$

Thus, for any $\varepsilon>0$ there exists $B \in \mathbb{R}^{|\sigma|, k}$ and $C \in \mathbb{R}^{|\tau|, k}$ with $k$ small (i.e. $k \ll|\sigma|$ and $k \ll|\tau|$ ) such that:

$$
\begin{equation*}
\left\|A_{\tau, \sigma}-B C^{t}\right\|_{F} \leq \varepsilon\left\|A_{\tau, \sigma}\right\|_{F} \tag{16}
\end{equation*}
$$

where $\|\cdot\|_{F}$ denotes the Frobenius norm, and we denote by $|\cdot|$ the cardinal of a set of indices. The storage requirement of the approximation is $k(|\tau|+|\sigma|)$, which is significantly lower than the $|\tau| .|\sigma|$ needed for the full storage of $A_{\tau, \sigma}$.

The $\mathcal{H}$-matrix format relies on:

- a block partition of the matrix that contains blocks satisfying the admissibility condition (14),
- a low-rank approximation of those admissible blocks with respect to a given precision (16).

In the following, we present the classical procedure for the approximation of an integral operator in the $\mathcal{H}$-matrix format [5].

### 2.2. Cluster tree and block tree partition

A good block partition of $A$ must contain a large number of admissible blocks [5]. Let $I=\{1, \ldots, n\}$ be the set of indices of the degrees of freedom of (13). The first step is the creation of a so-called cluster tree partition
$\mathcal{T}_{I}$ of $I$, defined in the following paragraph. Each node of that tree is a set of indices $\sigma \subset I$ that corresponds to a subdomain $\Omega_{\sigma}=\cup_{i \in \sigma} \operatorname{supp}\left(\phi_{i}\right)$ of $\Omega$, so that $\mathcal{T}_{I}$ equivalently define a partition tree of $\Omega$. The creation of $\mathcal{T}_{I}$ is so that the partition of $\Omega$ contains a large number of subdomains that are potentially admissible. The second step is the detection of admissible blocks : if $\sigma, \tau \in \mathcal{T}_{I}$ are so that $\Omega_{\sigma}$ and $\Omega_{\tau}$ satisfy (14), then those indices corresponds to a block of $A$. In practice, the detection of admissible blocks is done recursively in order to optimize the block partition of $A$. It results in a block tree partition presented in the next paragraph.

Remark 2.1. Only geometrical information is needed for the creation of the block partition.
Cluster tree partition. A tree $\mathcal{T}_{I}$, with nodes $T_{I}$, is called a cluster tree if the following conditions hold :
(1) $T_{I} \subset \mathcal{P}(I) \backslash\{\emptyset\}$, i.e. each node of $\mathcal{T}_{I}$ is a subset of the index set $I$,
(2) $I$ is the root of $\mathcal{T}_{I}$,
(3) If $\tau \in T_{I}$ is a leaf, then $|\tau| \leq C_{l e a f}$, i.e. the leaves consist of a relatively small number of indices,
(4) If $\tau \in T_{I}$ is not a leaf, then it has two sons and their union is disjoint.

The cluster tree $\mathcal{T}_{I}$ is recursively constructed with a function split. Starting from the root $\tau=I$ and from an initial tree $\mathcal{T}_{I}$ that contains only the root $I$, the algorithm proceeds as follow :

```
procedure }\mp@subsup{\mathcal{T}}{I}{}=\mathrm{ build_cluster_tree ( }\tau,\mp@subsup{\mathcal{T}}{I}{})\mathrm{ :
if }|\tau|\geq\mp@subsup{C}{leaf}{
    [\tau1, , _ ] = split (\tau),
    add \mp@subsup{\tau}{1}{}}\mathrm{ , and }\mp@subsup{\tau}{2}{}\mathrm{ in }\mp@subsup{\mathcal{T}}{I}{}\mathrm{ as sons of }\tau\mathrm{ ,
    call }\mp@subsup{\mathcal{T}}{I}{}=\mathrm{ build_cluster_tree ( }\mp@subsup{\tau}{1}{},\mp@subsup{\mathcal{T}}{I}{})\mathrm{ ,
    call }\mp@subsup{\mathcal{T}}{I}{}=\mathrm{ build_cluster_tree( }\mp@subsup{\tau}{2}{},\mp@subsup{\mathcal{T}}{I}{})\mathrm{ ,
end.
```

Remark 2.2. The value $C_{l e a f}=15$ classically leads to optimal computation time [5].
Note that the function $\left[\tau_{1}, \tau_{2}\right]=\operatorname{split}(\tau)$ typically uses geometrical information to split the node $\tau$. For each index $i \in \tau$, let $\mathbf{x}_{i}$ be the center of the support of $\phi_{i}$. For example, a well balanced tree can be obtained with a geometric bisection on the collection of points $\mathbf{x}_{i}$, meaning that if $\mathbf{x}_{i}$ is on one side of the corresponding hyperplane, $i$ will be assigned in $\tau_{1}$; otherwise in $\tau_{2}$. In figure 4 we illustrate the algorithm for the construction of the cluster tree partition.


Figure 4. Construction of the cluster partitioning of $\Omega$
Block tree partition. Given a cluster tree partition $\mathcal{T}_{I}$, the block tree partition $\mathcal{B}_{I, I}$ of $A$ is a tree that contains at each node a pair $(\tau, \sigma)$ of indices of $\mathcal{T}_{I}$. The leaves of $\mathcal{B}_{I, I}$ corresponds to admissible blocks (14). It can be constructed by the following procedure (see figure 5) initialized with $\tau=\sigma=I$ and $\mathcal{B}_{I, I}$ the block tree that contains only the root $(I, I)$ :
procedure $\mathcal{B}_{I, I}=$ build_block_tree $\left(\tau, \sigma, \mathcal{B}_{I, I}\right)$
if $(\tau, \sigma)$ is not admissible, and $|\tau| \geq C_{\text {leaf }}$, and $|\sigma| \geq C_{\text {leaf }}$
$S=\left\{\left(\tau^{*}, \sigma^{*}\right), \tau^{*}\right.$ son of $\tau, \sigma^{*}$ son of $\left.\sigma\right\}$,
add $S$ in $\mathcal{B}_{I, I}$ as son of $(\tau, \sigma)$,
for $\left(\tau^{*}, \sigma^{*}\right) \in S$


Figure 5. Construction of the block tree partition at level 1,2,3,4 of call of the build_block_tree function. The green blocks are admissible : at the end, only blocks of size $<C_{l e a f}$ are not admissible (pink). This matrix corresponds to a discretized 1D Laplace operator.

The complexity of algorithms for the creation of suitable cluster trees and block tree partitions has been analyzed in detail in [15]. For typical quasi-uniform grids, a "good" cluster tree can be created in $\mathcal{O}(n \log n)$ operations, the computation of the block partition can be accomplished in $\mathcal{O}(n)$ operations.

### 2.3. Approximation of the blocks

We continue this presentation by giving here some details on the construction of a low-rank approximation of each matrix block $A_{\tau, \sigma}$. The truncated singular value decomposition (SVD) is the optimal decomposition meaning that the relative precision $\varepsilon$ of (16) is archieved with minimal rank $k_{\tau, \sigma}$. But the SVD procedure requires the knowledge of the entire block $A_{\tau, \sigma}$.

To reduce the computational cost, the adaptive cross approximation (ACA) algorithm constructs a low rank approximation based on the knowledge of only a few particular rows and lines of the block. A first method for choosing these rows and lines, called ACA with full pivoting, has been proposed in [31]. In particular, it results in a quasi-optimal approximation. The drawback of this algorithm is that the determination of the ideal pivot requires the knowledge of the entire block, as described in the next paragraph. In [3] the authors proposed a different algorithm, called ACA with partial pivoting. In this case, the choice of the pivot is made in such a way that fewer block entries are needed.

ACA with full pivoting. We consider a $n$-by- $m$ matrix $M$. The adaptive cross approximation is a greedy procedure on the approximation $M^{k}=\sum_{\nu=1}^{k} a^{\nu} \otimes b^{\nu}$ of $M$. Each iteration consists in the following steps :
(1) Find the pivot $\left(i^{*}, j^{*}\right)$ such that:

$$
\begin{equation*}
\left(i^{*}, j^{*}\right)=\arg \max _{i j}\left|M_{i j}-M_{i j}^{k}\right| \tag{17}
\end{equation*}
$$

(2) Compute the two vectors $a_{i}^{k}=\left(M_{i j^{*}}-M_{i j^{*}}^{k}\right) /\left(M_{i^{*} j^{*}}-M_{i^{*} j^{*}}^{k}\right)$ and $b_{j}^{k}=\left(M_{i^{*} j}-M_{i^{*} j}^{k}\right)$
(3) Update the approximation $M^{k+1}=M^{k}+a^{k} \otimes b^{k}$.

The iterations are stopped when the desired precision is achieved using the condition (16).
In the general case there is no result on the rate of convergence. But when the matrix corresponds to a block of an discretized integral operator with asymptotically smooth the kernel $g$, it is proved that the convergence is exponential [5]. Moreover, this approximation is quasi-optimal. However the first step is to find the largest matrix entry (full pivoting (17)), which, as for the SVD, requires the knowledge of the full matrix $M$. The ACA algorithm with partial pivoting uses a different approach for the selection of the pivot $\left(i^{*}, j^{*}\right)$ which enables a significant reduction of the computational cost.

ACA with partial pivoting. The idea of partial pivoting is to maximise $\left|M_{i j}-M_{i j}^{k}\right|$ only for one of the two indices i or j and keep the other one fixed, i.e., we determine the maximal element in modulus in one particular row or one particular column. The new pivoting strategy consists at each iteration in:
(1) For a given index $i^{*}$, find the index $j^{*}$ such that:

$$
\begin{equation*}
j^{*}=\arg \max _{j}\left|M_{i^{*} j}-M_{i^{*} j}^{k}\right| \tag{18}
\end{equation*}
$$

(2) Compute the two vectors $a_{i}^{k}=\left(M_{i j^{*}}-M_{i j^{*}}^{k}\right) /\left(M_{i^{*} j^{*}}-M_{i^{*} j^{*}}^{k}\right)$ and $b_{j}^{k}=\left(M_{i^{*} j}-M_{i^{*} j}^{k}\right)$
(3) Update the approximation $M^{k+1}=M^{k}+a^{k} \otimes b^{k}$,
(4) Find the index $i^{*}$ such that:

$$
\begin{equation*}
i^{*}=\arg \max _{i}\left|M_{i j^{*}}-M_{i j^{*}}^{k}\right|=\arg \max _{i}\left|b_{j}^{k}\right| \tag{19}
\end{equation*}
$$

The stopping criterion (16) is replaced by a stagnation-based error estimator:

$$
\begin{equation*}
\left\|a^{k} \otimes b^{k}\right\|_{F} \leq \varepsilon\left\|M^{k}\right\|_{F} \tag{20}
\end{equation*}
$$

The particularity of this algorithm is that we do not have to compute all matrix entries of $M$. On the other hand, convergence is not guaranteed: one can find in [7] several counterexamples where ACA with partial pivoting is unable to reach the desired precision (16). Note that there exists a number of variants of this algorithm, such as the improved ACA and the Hybrid Cross Approximation [5, 7]: using additional heuristics, they try to improve on some typical failures of the basic ACA algorithm.

## 3. Numerical Results

### 3.1. Implementation

We present here some numerical results to illustrate the proposed method for the resolution of the corrector problem. For the implementation, we used the open-source boundary element library BEM $++[30]$ interfaced with the library Ahmed [4] for an efficient representation of the discretized integral operators by $\mathcal{H}$-matrices. All timings are reported for a laptop running the Python / $\mathrm{C}++$ code with a 2.7 GHz Intel Core i7 CPU. Three different microstructures are studied containing respectively 42, 98 and 203 randomly distributed inclusions, with nine different meshes ranging from 1344 to 89140 elements. The diffusion coefficient $\kappa$ takes the value 1 in the matrix and 100 in the inclusions. The precision used for the $\mathcal{H}$-matrix compression with ACA algorithm (see section 2 ) is set to $\varepsilon=1 e-3$.

Some graphical representations of the solution are shown in figure 6. In particular, we can observe the perturbation on the corrector field induced by the periodic copies of the inclusions where the diffusion coefficient is much higher than in the matrix.

### 3.2. Analysis of the $\mathcal{H}$-matrix approximation efficiency

We finally investigate here the scalability of the boundary element approach presented above. As a benchmark, we will also include results obtained with the same data, but replacing the periodic by the free-space Green's function. The behavior and scaling of the method with the number of degrees of freedom is well-known in this case, see e.g. [30].

We present in table 1 some data collected during the calculations. We observe that the memory storage for the $\mathcal{H}$-matrix representing the integral operator as well as the time used by the solver is slightly increased, but stays of the same order of magnitude in the periodic case in comparison to the benchmark free-space computation. By contrast, the assembly time through the ACA algorithm is much higher in the periodic case. This is due to

(A) Microstructure with 42 inclusions (meshed with 18424 triangles)

(в) Microstructure with 98 inclusions (43034 triangles)
(c) Microstructure with 203 inclusion (89140 triangles)

Figure 6. Representation of the reconstructed field of the corrector $u$, computed by the boundary element method for the three studied microstructures.
the much higher cost of evaluation of the periodic Green's function using the representation (8), by a factor of about 50 .

While no specific optimization effort has been made in the implementation of the periodic Green's function using formula (10), let us stress that a large part of the additional computational cost is due to the direct summation over the 27 free-space Green's function centered in the neighboring cells. In fact, due to the symmetry of the unit cell, only harmonics of even degree $l \geq 4$ and order $m$ multiple of 4 have a nonzero coefficient $\beta_{l}^{m}$ in (10) thus already limiting the impact of their (relatively expensive) evaluation in the total cost. For the degree $L=8$ used for the computations presented in this paper, only 11 solid spherical harmonics appear in the summation. To accelerate this expensive step in the computation when repeatedly assembling matrices for solving on a number of random geometries, one could employ an interpolation scheme between

| mesh size | Case | storage | compression | ACA time | solver |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 44784 triangles | Free-space | $615 \mathrm{Mb} / 15301 \mathrm{MB}$ | $4.02 \%$ | 63 s | 76 s |
| 44784 triangles | Periodic | $1012 \mathrm{Mb} / 15301 \mathrm{MB}$ | $6.6 \%$ | 3146 s | 118 s |
| 89140 triangles | Free-space | $1374 \mathrm{MB} / 60622 \mathrm{MB}$ | $2.26 \%$ | 128 s | 190 s |
| 89140 triangles | Periodic | $2210 \mathrm{MB} / 60622 \mathrm{MB}$ | $3.6 \%$ | 6402 s | 288 s |

Table 1. Data for the solution of the boundary element problem in a geometry with 203 inclusions
tabulated values of the periodic Green's function, precomputed using (10) and using the symmetry of the cell to reduce the memory requirements.

Figure 7 illustrates however that the scaling in performance with problem size is independent of the kernel and also of the number of inclusions in the computational geometry: in all cases we observe a memory and time cost scaling approximately as $\mathcal{O}\left(N^{1.3}\right)$.


Figure 7. Computational cost scaling as a function of mesh size for the periodic and the free-space cases.

## 4. Conclusion

We have proposed a new approach to the solution of the corrector problem by the use of boundary integral equations. The periodic boundary conditions are ensured through an approximation of the periodic Green's function. Then, the $\mathcal{H}$-matrix technique has been successfully applied for the resolution of the fully populated linear system araising from the Galerkin approximation of the corresponding boundary integral equation. The gain of memory and of computational costs provided by the $\mathcal{H}$-matrix format has allowed to deal with a very large number of degrees of freedom in the simulation. The first results are very promising, and we hope to further develop this method which will be presented in details and compared to existing approaches in [10].

## References

[1] G. Allaire. Homogenization and two-scale convergence. SIAM J. Math. Anal., 23(6):1482-1518, 1992.
[2] A. Barnett and L. Greengard. A new integral representation for quasi-periodic fields and its application to two-dimensional band structure calculations. Journal of Computational Physics, 229(19):6898-6914, 2010.
[3] M. Bebendorf. Approximation of boundary element matrices. Numerische Mathematik, pages 565-589, 2000.
[4] M. Bebendorf. Another software library on Hierarchical Matrices for elliptic differential equations (AHMED). Universität Leipzig, Fakultät für Mathematik und Informatik, 2005.
[5] M. Bebendorf. Hierarchical matrices: a means to efficiently solve elliptic boundary value problems, volume 63 of Lecture Notes in Computational Science and Engineering. Springer, 2008.
[6] A. Bensoussan, J.-L. Lions, and G. Papanicolaou. Asymptotic analysis for periodic structures, volume 5 of Studies in Mathematics and its Applications. North-Holland Publishing Co., Amsterdam, 1978.
[7] S. Börm, L. Grasedyck, and W. Hackbusch. Hierarchical matrices. Lecture notes, Max Planck Institut für Mathematik, 2004.
[8] J. M. Borwein, M. L. Glasser, R. C. McPhedran, J.G. Wan, and I. J. Zucker. Lattice Sums Then and Now, volume 150 of Encyclopedia of Mathematics and its Applications. Cambridge University Press, 2013.
[9] S. Brisard. Analyse morphologique et homogénéisation numérique, application à la pâte du ciment. PhD thesis, Université Paris-Est, 2011.
[10] P. Cazeaux and O. Zahm. Solving three-dimensional corrector problems in random homogenization with the boundary element method. in preparation, 2014.
[11] W. C. Chew, J. M. Jin, E. Michielssen, and J. Song. Fast and Efficient Algorithms in Computational Electromagnetics. Artech House, Boston, MA, 2001.
[12] J. W. Eischen and S. Torquato. Determining elastic behavior of composites by the boundary element method. Journal of applied physics, 74(1):159-170, 1993.
[13] P.P. Ewald. Die Berechnung optischer und elektrostatischer Gitterpotentiale. Ann. Physik., 369(3):253-287, 1921.
[14] A. Gloria and F. Otto. An optimal error estimate in stochastic homogenization of discrete elliptic equations. The Annals of Applied Probability, 22(1):1-28, 2012.
[15] L. Grasedyck. Theorie und Anwendungen Hierarchischer Matrizen. PhD thesis, University of Kiel, 2001.
[16] L. Greengard and J. Helsing. On the numerical evaluation of elastostatic fields in locally isotropic two-dimensional composites. Journal of the Mechanics and Physics of Solids, 46(8):1441-1462, 1998.
[17] L. Greengard and M. C. Kropinski. Integral equation methods for stokes flow in doubly-periodic domains. Journal of engineering mathematics, 48(2):157-170, 2004.
[18] L. Greengard and V. Rokhlin. A fast algorithm for particle simulations. Journal of computational physics, 73(2):325-348, 1987.
[19] W. Hackbusch. A Sparse Matrix Arithmetic Based on H-Matrices. Part I: Introduction to H-Matrices. Computing, 62(2):89-108, April 1999.
[20] W. Hackbusch and L. Grasedyck. An introduction to hierarchical matrices. Mathematica Bohemica, 127(2):229-241, 2002.
[21] Z. Hashin and S. Shtrikman. A variational approach to the theory of the elastic behaviour of multiphase materials. Journal of the Mechanics and Physics of Solids, 11(2):127-140, 1963.
[22] K Houzaki, N. Nishimura, and Y. Otani. An fmm for periodic rigid-inclusion problems and its application to homogenisation. Contemporary mathematics, 408:81-98, 2006.
[23] G. C. Hsiao and W. L. Wendland. Boundary integral equations, volume 164 of Applied Mathematical Sciences. Springer-Verlag, Berlin, 2008.
[24] S.M. Kozlov. Averaging of random operators. Mathematics of the USSR-Sbornik, 37(2):167, 1980.
[25] Y. J. Liu, S. Mukherjee, N. Nishimura, M. Schanz, W. Ye, A. Sutradhar, E. Pan, N. A. Dumont, A. Frangi, and A. Saez. Recent advances and emerging applications of the boundary element method. Applied Mechanics Review, 64(3), 2011.
[26] V. P. Nguyen, M. Stroeven, and L. J. Sluys. Multiscale continuous and discontinuous modeling of heterogeneous materials: a review on recent developments. Journal of Multiscale Modelling, 3(4):1-42, 2011.
[27] Y. Otani and N. Nishimura. A fast multipole boundary integral equation method for periodic boundary value problems in three-dimensional elastostatics and its application to homogenization. International Journal for Multiscale Computational Engineering, 4(4), 2006.
[28] G. J. Rodin and J. R. Overfelt. Periodic conduction problems: the fast multipole method and convergence of integral equations and lattice sums. Proceedings of the Royal Society of London. Series A: Mathematical, Physical and Engineering Sciences, 460(2050):2883-2902, 2004.
[29] E. Sanchez-Palencia. Homogenization in mechanics. a survey of solved and open problems. Rend. Sem. Mat. Univers. Politecn. Torino, 44(1), 1986.
[30] W. Smigaj, S. Arridge, T. Betcke, J. Phillips, and M. Schweiger. Solving Boundary Integral Problems with BEM ++ . preprint, 2012.
[31] E. E. Tyrtyshnikov. Incomplete cross approximation in the mosaic-skeleton method. Computing, 64(4):367-380, 2000.
[32] J. Yuan, Y. Y. Lu, and X. Antoine. Modeling photonic crystals by boundary integral equations and Dirichlet-to-Neumann maps. Journal of Computional Physics, 227(9):4617-4629, 2008.


[^0]:    * Thanks to A. Nouy, F. Legoll, Y. Assami, A. Obliger, W. Minvielle for fruitful discussions.
    ** The authors would like to CEMRACS organizers and participants for a great stay.
    *** This work was financially supported by EDF R\&D.
    1 MATHICSE, Chair of Computational Mathematics and Simulation Sciences, Ecole Polytechnique Fédérale de Lausanne, Station 8, CH-1015 Lausanne, Szitwerland. (paul.cazeaux@epfl.ch)
    ${ }^{2}$ GeM, UMR CNRS 6183, École Centrale de Nantes, 1 rue de la Noë, BP92101, 44321 Nantes Cedex 3, France. (olivier.zahm@ecnantes.fr)

