

Quasi-Newton Methods

BFGS (Broyden, Fletcher, Goldfarb, Shanno)

Consider (again) $m_k(p) = F_k + \nabla F_k^T p + \frac{1}{2} p^T B_k p$

for B_k SPD and B_k approx. Hessian

Proceeding as for Newton: $p_k = -B_k^{-1} \nabla F_k$

$$x_{k+1} = x_k + \alpha_k p_k$$

where α_k satisfies Wolfe cond.s

In the next step we want new model

$$m_{k+1}(p) = F_{k+1} + \nabla F_{k+1}^T p + \frac{1}{2} p^T B_{k+1} p$$

s.t. $\nabla m_{k+1}(0) = \nabla F_{k+1}$ (obvious)

$$\nabla m_{k+1}(-\alpha_k p_k) = \nabla F_k \quad \{ \rightarrow$$

$$\nabla F_{k+1} - \alpha_k B_{k+1} p_k = \nabla F_k \quad \Leftrightarrow$$

$$B_{k+1} \alpha_k p_k = \nabla F_{k+1} - \nabla F_k$$

Set $x_{k+1} - x_k =: s_k$ $y_k := \nabla F_{k+1} - \nabla F_k$:

$$B_{k+1} s_k = y_k \quad (\text{second eq. / cond.})$$

$$B_{k+1} \text{ SPD} \rightarrow s_k^T B_{k+1} s_k = s_k^T y_k > 0$$

poses constraints on L.S. Wolfe or SW/cond.s

guarantee this curvature cond. $\#$

$$\text{curv. cond: } \nabla F_{k+1}^T p_k \geq c_2 \nabla F_k^T p_k \Rightarrow (s_k = \alpha_k p_k)$$

$$\nabla F_{k+1}^T s_k \geq c_2 \nabla F_k^T s_k$$

$$(\nabla P_{k+1}^* - \nabla P_k^*)^T s_k \geq (c_2 - 1) \nabla P_k^* s_k > 0 \quad / \quad 106$$

since $c_2 < 1$ and $\nabla P_k^* p_k < 0$ ($\alpha_k > 0$)

Many possible SPD choices B_{k+1}

For BFGS:

$$\text{let } H_k = B_k^{-1} \quad (\text{also SPD}) \rightarrow H_{k+1} y_k = s_k$$

$$\text{and } \min_H \|H - H_k\| \rightarrow H_{k+1}$$

$$\text{subject to } H^T = H \text{ and } H y_k = s_k$$

$$\text{Take } \|\bar{G}_k^{-1/2} (H - H_k) \bar{G}_k^{-1/2}\|_F \text{ min where}$$

$$\bar{G}_k \text{ average Hessian } \bar{G}_k = \int_0^1 \nabla^2 P(x_k + \tau \alpha_k p_k) d\tau$$

$$\text{(BFGS)} \quad \text{Then } H_{k+1} = (I - p_k s_k^T) H_k (I - p_k y_k^T) + p_k s_k s_k^T$$

$$\text{where } p_k = (y_k^T s_k)^{-1}$$

How to pick H_0 ?

exact Hessian or good approx. at x_0
(e.g. by finite diff's)

constant times ∇ (poss. 1)

(p. 143) Use $\tilde{H}_0 = I$ to get p_k , compute y_k, s_k

and then set H_0 to $\frac{y_0^T s_0}{y_0^T y_0} I$ before

computing the update to get H_1

The BFGS update of B_k (corr. to update H_k)

$$B_{k+1} = B_k - \frac{B_k s_k s_k^T B_k}{s_k^T B_k s_k} + \frac{y_k y_k^T}{y_k^T s_k}$$

(However, requires Chol. decomp $\rightarrow O(n^3)$ cost)
 poss. to update Chol. factors directly $\rightarrow O(n^2)$ cost.

If B_k SPD then B_{k+1} SPD (assuming d_k following WC or SWC)

$$z^T H_{k+1} z = \frac{z^T H_k z}{z^T s_k} \quad (z \neq 0)$$

$$\frac{z^T (\mathbb{I} - \rho_k s_k y_k^T) H_k (\mathbb{I} - \rho_k y_k s_k^T) z}{w^T} + \rho_k \frac{z^T s_k s_k^T z}{z^T s_k} =$$

$$\frac{w^T H_k w}{w^T} + (y_k^T s_k)^{-1} (z^T s_k)^2$$

> 0 if $w \neq 0$

> 0 if $z^T s_k \neq 0$

if $z^T s_k = 0$ then $w \neq 0$ and $z^T H_{k+1} z > 0$

if $z^T s_k = 0$ then $z^T H_{k+1} z = z^T H_k z > 0$

An alternative for BFGS is SR1 method that satisfies some remarkable properties (later), and often approx. Hessian better than BFGS.

SR1 req. only rank 1 update. Since it may not maintain pos. def. better used with TR than with LS. (see Alg. 6.2)

$$B_{k+1} = B_k + \sigma v v^T \quad \sigma = \pm 1$$

$$y_k = (B_k + \sigma v v^T) s_k = B_k s_k + \sigma v v^T s_k \quad (\Leftrightarrow)$$

← scalar

$$y_k - B_k s_k = v (\sigma v^T s_k)$$

Hence $v = \delta (y_k - B_k s_k) \rightarrow$

$$(y_k - B_k s_k) = \sigma \delta^2 (s_k^T (y_k - B_k s_k)) (y_k - B_k s_k)$$

If $y_k - B_k s_k = 0 \rightarrow v = 0 : B_{k+1} = B_k$

If not, in principle: $\begin{cases} \delta = \frac{1}{\|s_k\|} |s_k^T (y_k - B_k s_k)|^{-\frac{1}{2}} \\ \sigma = \text{sign}(s_k^T (y_k - B_k s_k)) \end{cases}$

if δ small problem $\epsilon \rightarrow$

if $|s_k^T (y_k - B_k s_k)| \geq r \|s_k\| \|y_k - B_k s_k\|$

for a small $r \in (0, 1)$, e.g., $r = 10^{-8}$ (book)

then update is done, otherwise $B_{k+1} = B_k$

Theo 6.1: Suppose $f: \mathbb{R}^n \rightarrow \mathbb{R}$ defined as

$$f(x) = \frac{1}{2} x^T A x + b^T x, \text{ where } A \text{ SPD. Then for}$$

any starting point x_0 and any symmetric starting matrix H_0 , the iterates $\{x_k\}$ from SR1 with

$$\begin{cases} H_{k+1} = H_k + \frac{(s_k - H_k y_k)(s_k - H_k y_k)^T}{(s_k - H_k y_k)^T y_k} \\ p_k = -H_k \nabla f_k, \quad x_{k+1} = x_k + p_k \end{cases}$$

converge to x^* in n steps, provided

$$(s_k - H_k y_k)^T y_k \neq 0 \text{ for all } k.$$

Moreover, if n steps are performed and the search directions p_i linearly indep, then

$$H_n = A^{-1}.$$

Proof: (see book, simple by induction showing

$$H_{k+1} y_j = s_j \quad j = 0 \dots k)$$

Theorem 6.2: Suppose f twice cont. diff. and Hessian bnd and Lip. cont. in nbhd of point x^* .

Let $\{x_k\}$ be any seq. of iterates s.t. $x_k \rightarrow x^*$ for some $x^* \in \mathbb{R}^n$. In addition, suppose for all k

$$|s_k^T (y_k - B_k s_k)| \geq r \|s_k\| \|y_k - B_k s_k\| \text{ for some } r \in (0, 1)$$

and that the s_k are uniformly indep (Form space of max dim). Then $\lim_{k \rightarrow \infty} \|B_k - \nabla^2 f(x^*)\| = 0$

If f twice cont. diff. and the level set $R = \{x \mid f(x) \leq f(x_0)\}$ is convex, and Hessian $G(x)$ satisfies $0 < m \leq z^T G(x) z \leq M < \infty$ for all $x \in R$ and all $z \in \mathbb{R}^n$. In addition if B_0 SPD. Then $\{x_k\}$ generated by BFGS

implemented by Alg. 6.1 converges to the minimizer x^* of f .

Note that assumptions are significantly stronger than usual.

Theo.
If f twice cont. diff. and $\{x_k\}$ generated by

BFGS converge to x^* s.t. $\sum_{k=0}^{\infty} \|x_k - x^*\| < \infty$

and Hessian $G(x)$ lip. cont. at x^* ,

$$\|G(x) - G(x^*)\| \stackrel{\text{(locally)}}{\leq} L \|x - x^*\| \text{ for all } x \text{ near } x^*$$

($L > 0$), then $x_k \rightarrow x^*$ superlinearly.

~~Theo~~ Proof, see book.

Large-scale Unconstrained Optimization

Inexact Newton Methods

Solve $\nabla^2 f_k p_k = -\nabla f_k$ or $B_k p_k = -\nabla f_k$

use CG for approx. sol.

residual: $r_k = \nabla^2 f_k p_k + \nabla f_k$

termination cond. $\|r_k\| \leq \eta_k \|\nabla f_k\|$

for seq. $\{\eta_k\}$ w. $0 < \eta_k < 1$ ($\forall k$)

\hookrightarrow forcing seq.

Assume $\nabla^2 f(x)$ cont in nbhood men. x^*
and pos. def. at x^*

Use a forcing seq. $\eta_k \leq \eta$ $\eta \in (0, 1)$

For x suff. close to x^* , $\|(\nabla^2 f(x))^{-1}\| \leq L$

$$p_k = (\nabla^2 f_k)^{-1} (r_k + \nabla f_k)$$

$$\|p_k\| \leq L (\|\nabla f_k\| + \|r_k\|) \leq 2L \|\nabla f_k\|$$

$\leq \eta \|\nabla f_k\|$

~~$$\nabla f_{k+h} = \nabla f(x_k + \alpha p_k) = \nabla f_k + \nabla^2 f_k p_k + \dots$$~~

$$\nabla f_{k+h} = \nabla f(x_k + \alpha p_k) =$$

$$\nabla f_k + \underbrace{\nabla^2 f_k p_k}_{r_k - \nabla f_k} + \int_0^1 [\nabla^2 f(x_k + t p_k) - \nabla^2 f(x_k)] p_k dt$$

$$\| \dots \| = o(\|p_k\|)$$

$$= r_k + o(\|\nabla f_k\|)$$

$$\begin{aligned} \|\nabla f_{k+1}\| &\leq \eta_k \|\nabla f_k\| + o(\|\nabla f_k\|) \\ &\leq (\eta_k + o(1)) \|\nabla f_k\| \quad (\text{too many braces in b/w}) \end{aligned}$$

$o(1)$ term $\rightarrow 0$, so assume $o(\|\nabla f_k\|)$ term smaller than $(1-\eta)/2$

$$\|\nabla f_{k+1}\| \leq \frac{1+\eta}{2} \|\nabla f_k\| \quad (\text{close enough to } x^*)$$

Start close enough to $x^* \rightarrow$ have this each step.

$$\nabla f_k = \nabla^2 f(x^*)(x_k - x^*) + o(\|x_k - x^*\|)$$

$$\|\nabla f_k\| \leq \|\nabla^2 f(x^*)\| \|e_k\| + o(\|e_k\|)$$

Above $\frac{\|\nabla f_{k+1}\|}{\|\nabla f_k\|} \leq \eta_k + o(1)$ or

$$\left(\frac{\|\nabla f_{k+1}\|}{\|\nabla f_k\|} \leq \frac{1+\eta}{2} \right)$$

So, for $\eta_k \rightarrow 0$ we get

$$\lim_{k \rightarrow \infty} \frac{\|\nabla f_{k+1}\|}{\|\nabla f_k\|} = 0 \quad \text{superlinear conv.}$$

of grad. to zero (implies superlin conv x_k to x^*)
 \rightarrow SPD $\nabla^2 f$ implies $\nabla f = 0 \Leftrightarrow x_k = x^*$

If $\nabla^2 f$ lip. cont. near x^* , then

$$\|\nabla f_{k+1}\| = O(\|\nabla f_k\|^2) \quad (\text{quadr. conv.})$$

$$\eta_k = O(\|\nabla f_k\|)$$

$\nabla^2 f$ lip cont in nbhood x^* and
pos. def x^*

Forcing seq. $\{\eta_k\}$ with $\eta_k < \eta \in (0, 1)$

For x suff. close to x^* $\|\nabla^2 f(x)^{-1}\| \leq L$

$$p_k = (\nabla^2 f_k)^{-1} (r_k + \nabla f_k) \quad \|r_k\| \leq \eta_k \|\nabla f_k\|$$

$$\|p_k\| \leq L(\|\nabla f_k\| + \|r_k\|)$$

$$\leq 2L \|\nabla f_k\|$$

$$\nabla f_{k+1} = \nabla f(x_k + \alpha_k p_k) =$$

$$\nabla f_k + \nabla^2 f_k p_k + \underbrace{\int_0^1 [\nabla^2 f(x_k + t p_k) - \nabla^2 f(x_k)] p_k dt}_{\| \cdot \| \leq \frac{1}{2} \gamma \|p_k\|^2}$$

$$\begin{aligned} \|\nabla^2 f(x_k + t p_k) - \nabla^2 f(x_k)\| &\leq \\ \gamma \|t p_k\| &= \\ \gamma t \|p_k\| & \end{aligned}$$

These results \rightarrow Theo. 7.2

$$\text{sugg. for } \eta_k: \begin{cases} \eta_k = \min(1/2, \|\nabla f_k\|^{-1/2}) \\ \eta_k = \min(1/2, \|\nabla f_k\|) \end{cases}$$

Line Search Newton CG for

$$\mathbf{B}_k \mathbf{p} = -\nabla f_k$$

Choose x_0 ; $k=0$; tol; comp $\nabla f(x_0)$
while $\|\nabla f_k\| \leq > \text{tol}$,

$$\varepsilon_k = \eta_k \|\nabla f_k\|$$

$$z_0 = 0; r_0 = +\nabla f_k; d_0 = -r_0; j=0$$

while $\|r_j\| > \varepsilon_k$,

if $d_j^T \mathbf{B}_k d_j \leq 0$ (neg. curv)
if $j=0$ then $p_k = -\nabla f_k$
else $p_k = z_j$

$$\alpha_j = r_j^T r_j / d_j^T \mathbf{B}_k d_j$$

$$z_{j+1} = z_j + \alpha_j d_j$$

$$r_{j+1} = r_j + \alpha_j \mathbf{B}_k d_j$$

if $\|r_{j+1}\| > \varepsilon_k$

$$\beta_{j+1} = r_{j+1}^T r_{j+1} / r_j^T r_j$$

$$d_{j+1} = -r_{j+1} + \beta_{j+1} d_j$$

end while

$$x_{k+1} = x_k + \alpha_k p_k \quad (\text{where } \alpha_k \text{ sat.})$$

Wolfe/Armijo cond. and $\alpha_k = 1$
when possible \rightarrow backtracking)

end

Essentially CG alg. with add
test for B_k indef and special
update w. direction of negative
curvature.

Choice of η_k as suggested above.

For SPD $B_k = \nabla^2 f_k$ $z_j \rightarrow p_k^N$

!

B_k needed only for matrix-vector
product. So, matrix-free impl. poss.

$$\nabla^2 f_k \cdot d \approx \frac{\nabla^2 f(x_k + hd) - \nabla^2 f_k}{h} \quad \leftarrow \text{known}$$

For small but not too small h ;

essentially strike balance between
accuracy in exact arithmetic $h=0$
and cancellation (numerical error)

if $\frac{\|\nabla^2 f(x_k + hd) - \nabla^2 f_k\|}{\|\nabla^2 f_k\|}$ too small.

If B_k too ill-cond. LS-CG not so
effective.

Alg. 7.2

p. 171

Trust-Region CG

Use modified CG to solve/minimize approximately for ∇R step p

$$m_k(p) = f_k + g_k^T p + \frac{1}{2} p^T B_k p$$

(B may be indefinite) $\|p\| \leq \Delta$

* Use conv. criterion similar to $\frac{1}{2}$

CG-LS \rightarrow forcing seq. η_k (ϵ_k)

* Stop on direction of negative curv. find min on bound ∇R

$$* \begin{aligned} z_0 &= 0 \\ z_1 &= - \frac{\nabla f_k^T \nabla f_k}{\nabla f_k^T B_k \nabla f_k} \nabla f_k \end{aligned}$$

if $\|z_1\| \leq \Delta \rightarrow$ Cauchy pt.

subseq. steps improve

if $\|z_1\| > \Delta \rightarrow$ Cauchy pt and stop

* ~~if~~ $\|z_j\| > \|z_{j-1}\| > \dots$

So, can always stop at ∇R bound.

Prop. follows from prop. of CG and choice $z_0 = 0$.

Theo 7.3 $\{z_j\}$ from Alg 7.2 satisfies

$$0 = \|z_0\| < \|z_1\| < \dots < \|p_k\| \leq \Delta_k$$

The CG iter. can be pre conditioned to accelerate convergence.
Often used are variants of Inexact Modified Cholesky.

The choice of step when neg. curv. is discovered, can be poor.

As alternative we can use Lanczos method for tridiag. reduction/proj. of B_k (as in CG) and then use this directly for minimization.

$$B_k p = -\nabla f_k$$

$$* \rightarrow q_j = \pm r_j / \|r_j\|$$

$$r_0 = \nabla f_k$$

Iterate as for CG ($z_0 = 0$) =

$$B_k Q_j = Q_j T_j + \beta_j q_{j+1} e_j^T$$

$$T_j = Q_j^T B_k Q_j \quad \text{and} \quad Q_j^T Q_j = I^*$$

Find minimum for ∇R problem from $R(Q_j) = \rightarrow p = Q_j w$

$$\min_{w \in \mathbb{R}^j} f_k + \nabla f_k^T Q_j w + \frac{1}{2} w^T Q_j^T B_k Q_j w$$

$$\text{subj } \|Q_j w\| = \Delta$$

Simplify: $q_1 = \nabla f_k / \|\nabla f_k\|$ and $Q_j^T Q_j = I$

$$\rightarrow \nabla f_k^T Q_j w = \|\nabla f_k\| e_1^T w$$

$$Q_j^T B_k Q_j = T_j \quad \text{tridiag.}$$

$$\|Q_j w\| = \Delta \Rightarrow \|w\| = \Delta$$

$$\min_{w \in \mathbb{R}^j} f_k + \|\nabla f_k\| e_1^T w + \frac{1}{2} w^T T_j w$$

Since T_j is tridiag and small

we can easily factorize

$\nabla_j + dI$ and follow the (nearly) exact TR model/subproblem solution approach of section 4.3.

Limited-memory BFGS

Inverse updated BFGS =

$$x_{k+1} = x_k - \alpha_k H_k \nabla f_k$$

$$H_{k+1} = V_k^T H_k V_k + \rho_k s_k s_k^T \quad \text{where}$$

$$\rho_k = \frac{1}{y_k^T s_k}, \quad V_k = I - \rho_k y_k s_k^T$$

$$s_k = x_{k+1} - x_k, \quad y_k = \nabla f_{k+1} - \nabla f_k$$

H_k in general dense \rightarrow expensive for large problems.

Solution: store H_k implicitly (approx.) by keeping m pairs $\{s_i, y_i\}$

$H_k \nabla f_k$ can then be computed efficiently

Keep m pairs $\{s_i, y_i\} \quad i = k-m, \dots, k-1$

H_k^0 given (simple SPD matrix)

$$\begin{aligned} H_k &= V_{k-1}^T H_k^0 V_{k-1} + \rho_{k-1} s_{k-1} s_{k-1}^T = \\ &= V_{k-1}^T \left(V_{k-2}^T H_k^0 V_{k-2} + \rho_{k-2} s_{k-2} s_{k-2}^T \right) V_{k-1} \\ &\quad + \rho_{k-1} s_{k-1} s_{k-1}^T \end{aligned}$$

$$H_k = H_k^{(k)}$$

$$V_k = I - P_k Y_k S_k^T$$

$$= V_{k-1}^T V_{k-2}^T \dots V_{k-m}^T H_k^{(0)} V_{k-m} \dots V_{k-2} V_{k-1} \\ + P_{k-m} (V_{k-1}^T - V_{k-m+1}^T) S_{k-m} S_{k-m}^T V_{k-m+1} \dots V_{k-1} \\ + \vdots \\ + P_{k-1} S_{k-1} S_{k-1}^T$$

↓ Can be implemented eff. using recursion.
→ Alg. 7.4 (explicitly unrolled)

Poss. choice $H_k^{(0)} = \gamma_k I$ where

$$\gamma_k = \frac{S_{k-1}^T Y_{k-1}}{Y_{k-1}^T Y_{k-1}}$$

L-BFGS

Choose $x_0, m > 0$

$k=0$

while $\|\nabla f_k\| > \text{tol}$ and $k < \text{maxit}$

Choose $H_k^{(0)}$

$$P_k = -H_k^{(0)} \nabla f_k \quad (\text{using approx repr})$$

$$x_{k+1} = x_k + \alpha_k P_k \quad (\alpha_k \text{ sat's Wolfe cond.})$$

if $k > m$

Discard $\{S_{k-m}, Y_{k-m}\}$

$$\text{Compute and save } S_k = x_{k+1} - x_k \\ Y_k = \nabla f_{k+1} - \nabla f_k$$

$k = k+1$

end

$H_k^{(0)}$ may vary
per step

Compact Representation of BFGS updating

For TR we need limited memory BFGS representation for B_k itself.

*1

$$B_{k+1} = B_k - \frac{B_k s_k s_k^T B_k}{s_k^T B_k s_k} + \frac{y_k y_k^T}{y_k^T s_k}$$

$$\left(B_{k+1} s_k = y_k \rightarrow B_k s_{k-1} = y_{k-1} \right.$$

$$B_{k+1} = B_k \left(I - \frac{s_k s_k^T}{s_k^T B_k s_k} \right) + \frac{y_k y_k^T}{y_k^T s_k}$$

$$= \left(B_{k-1} \left(I - \frac{s_{k-1} s_{k-1}^T}{s_{k-1}^T B_{k-1} s_{k-1}} \right) B_{k-1} \right)$$

$$= \left(B_{k-1} \left(I - \frac{s_{k-1} s_{k-1}^T}{s_{k-1}^T B_{k-1} s_{k-1}} \right) B_{k-1} + \frac{y_{k-1} y_{k-1}^T}{y_{k-1}^T s_{k-1}} \right) \left(I - \dots \right) \left(\right)$$

$$+ \frac{y_k y_k^T}{y_k^T s_k} \rightarrow$$

$$B_k = B_0 - \begin{bmatrix} B_0 s_k & y_k \end{bmatrix} \begin{bmatrix} s_k^T B_0 s_k & L_k \\ L_k^T & -D_k \end{bmatrix}^{-1} \begin{bmatrix} s_k^T B_0 \\ y_k^T \end{bmatrix}$$

*2

$$S_k = [s_0 \dots s_{k-1}] \quad Y_k = [y_0 \dots y_{k-1}]$$

$$(L_k)_{ij} = \begin{cases} s_{i-1}^T y_{j-1} & \text{if } i > j \\ 0 & \text{otherwise} \end{cases}$$

$$D_k = \text{diag}(s_0^T y_0, \dots, s_{k-1}^T y_{k-1})$$

If B_0 SPD and $\{s_i, y_i\}_{i=0}^{k-1}$ sat. S

$s_i^T y_i > 0$ for all i then B_k (*1) =

B_k (*2). and $\begin{bmatrix} x & \vdots \\ \vdots & \ddots \end{bmatrix}^{-1}$ exists.

(note inverse exists because B_0 SPD and assumed $s_0^T y_0 > 0$)

$$\begin{aligned}
 B_1 &= B_0 - [B_0 s_0 \ y_0] \begin{bmatrix} s_0^T B_0 s_0 & \cancel{0} \\ \cancel{0} & -s_0^T y_0 \end{bmatrix}^{-1} \begin{bmatrix} s_0^T B_0 \\ y_0^T \end{bmatrix} \\
 &= B_0 - [B_0 s_0 \ y_0] \begin{bmatrix} (s_0^T B_0 s_0)^{-1} s_0^T B_0 \\ -(s_0^T y_0)^{-1} y_0^T \end{bmatrix} \\
 &= B_0 - \begin{bmatrix} B_0 s_0 s_0^T B_0 / (s_0^T B_0 s_0)^{-1} & -y_0 y_0^T (s_0^T y_0)^{-1} \end{bmatrix}
 \end{aligned}$$

(verified for $k=1 \rightarrow B_1$)

Assume true for $k = 1 \dots m-1$

~~$$B_{m-1} = B_0 - [B_0 s_{m-1} \ y_{m-1}] \begin{bmatrix} s_{m-1}^T B_{m-1} s_{m-1} & L_{m-1} \\ L_{m-1}^T & -D_{m-1} \end{bmatrix}^{-1} \begin{bmatrix} s_{m-1}^T B_0 \\ y_{m-1}^T \end{bmatrix}$$~~

$$B_{m-1} = B_0 - [B_0 s_{m-1} \ y_{m-1}] \begin{bmatrix} s_{m-1}^T B_{m-1} s_{m-1} & L_{m-1} \\ L_{m-1}^T & -D_{m-1} \end{bmatrix}^{-1} \begin{bmatrix} s_{m-1}^T B_0 \\ y_{m-1}^T \end{bmatrix}$$

$$B_m =$$

For $k > m$ i.f.s, keeping only m pairs of vectors:

$$B_0 = \delta_k \mathbf{I} = \gamma_k^{-1} \mathbf{I} \quad (\gamma_k = \frac{\mathbf{s}_{k-1}^T \mathbf{y}_{k-1}}{\mathbf{y}_{k-1}^T \mathbf{y}_{k-1}})$$

$$\mathbf{S}_k = [\mathbf{s}_{k-m} \ \mathbf{s}_{k-m+1} \ \dots \ \mathbf{s}_{k-1}] \quad \mathbf{Y}_k = [\mathbf{y}_{k-m} \ \dots \ \mathbf{y}_{k-1}]$$

$$B_k = \delta_k \mathbf{I} - [\delta_k \mathbf{S}_k \ \mathbf{Y}_k] \begin{bmatrix} \delta_k \mathbf{S}_k^T \mathbf{S}_k & \mathbf{L}_k \\ \mathbf{L}_k^T & -\mathbf{D}_k \end{bmatrix}^{-1} \begin{bmatrix} \delta_k \mathbf{S}_k^T \\ \mathbf{Y}_k^T \end{bmatrix}$$

$$(L_k)_{ij} = \begin{cases} \mathbf{s}_{k-m-1+i}^T \mathbf{y}_{k-m-1+j} & \text{if } i > j \\ 0 & \text{otherwise} \end{cases}$$

$$\mathbf{D}_k = \text{diag}(\mathbf{s}_{k-m}^T \mathbf{y}_{k-m} \ \dots \ \mathbf{s}_{k-1}^T \mathbf{y}_{k-1})$$

~~$$B_k = \mathbf{L} \mathbf{S}^T - \mathbf{W} \mathbf{F}^{-1} \mathbf{W}^T$$~~

$$B_k = \delta \mathbf{I} - \mathbf{W} \mathbf{F}^{-1} \mathbf{W}^T$$

$$\mathbf{F} = \begin{bmatrix} \delta \mathbf{S}^T \mathbf{S} & \mathbf{L} \\ \mathbf{L}^T & -\mathbf{D} \end{bmatrix} = \begin{bmatrix} \mathbf{S}^T \mathbf{S} & \mathbf{L} \\ \mathbf{L}^T & -\mathbf{D} \end{bmatrix}$$

$$\begin{bmatrix} \delta^{1/2} (\mathbf{S}^T \mathbf{S})^{1/2} & \mathbf{0} \\ \delta^{-1/2} \mathbf{L}^T (\mathbf{S}^T \mathbf{S})^{-1/2} & -\mathbf{G}^{1/2} \end{bmatrix} \begin{bmatrix} \delta^{1/2} (\mathbf{S}^T \mathbf{S})^{1/2} \delta^{-1/2} (\mathbf{S}^T \mathbf{S})^{-1/2} \mathbf{L} \\ \mathbf{0} & \mathbf{G}^{1/2} \end{bmatrix}$$

$\mathbf{L} \qquad \mathbf{R}$

$$\mathbf{G} = \delta \mathbf{L}^T (\mathbf{S}^T \mathbf{S})^{-1} \mathbf{L} + \mathbf{D}$$

$$B_k = \delta \mathbf{I} - (\mathbf{W} \mathbf{R}^{-1}) (\mathbf{L}^T \mathbf{W}^T) = \delta \mathbf{I} - \mathbf{U} \mathbf{V}^T$$

$$B_k^{-1} = \delta^{-1} \mathbf{I} + \mathbf{U} \mathbf{Z} \mathbf{V}^T \quad \rightarrow$$

$$\mathbf{I} - \delta^{-1} \mathbf{U} \mathbf{V}^T + \delta \mathbf{U} \mathbf{Z} \mathbf{V}^T - \mathbf{U} \mathbf{Z} \mathbf{V}^T \mathbf{U} \mathbf{V}^T = \mathbf{I}$$

$$\mathbf{U} (-\delta^{-1} \mathbf{I} + \delta \mathbf{Z} - \mathbf{Z} \mathbf{V}^T \mathbf{U}) \mathbf{V}^T = \mathbf{0}$$

$$-\delta^{-1} \mathbf{I} + \delta \mathbf{Z} - \mathbf{Z} \mathbf{V}^T \mathbf{U} = \mathbf{0} \Leftrightarrow -\mathbf{I} + \delta^2 \mathbf{Z} - \delta \mathbf{Z} \mathbf{V}^T \mathbf{U} = \mathbf{0}$$

$$\mathbf{Z} (\delta^2 \mathbf{I} - \delta \mathbf{V}^T \mathbf{U}) = \mathbf{I} \Leftrightarrow \mathbf{Z} = (\delta^2 \mathbf{I} - \delta \mathbf{V}^T \mathbf{U})^{-1}$$

\hookrightarrow ek rek

$$\delta_k \mathbf{I} + \begin{bmatrix} \blacksquare & \blacksquare \\ \blacksquare & \blacksquare \end{bmatrix}$$

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So, using ~~matrix~~ this easy to compute inverse for B_k , we can easily implement Dogleg or other methods for TR (approx) solution.

$$\text{Minimize } \frac{1}{2} \sum_i r_i^2(x) = f(x)$$

$$\begin{aligned} \frac{\partial f}{\partial x_k} &= \sum_i r_i(x) \cdot \frac{\partial r_i}{\partial x_k} \\ &= r^T \left(\frac{\partial r}{\partial x} \right)_k \\ &= (y^T r)_k \quad \begin{matrix} (k^{\text{th}} \text{ column}) \\ k^{\text{th}} \text{ coeff.} \end{matrix} \end{aligned}$$

$$g_{ij} = \left(\frac{\partial r_i}{\partial x_j} \right)$$

$$\nabla f = g^T r$$

$$\begin{aligned} H_{jk} = H_{kj} &= \frac{\partial^2 f}{\partial x_j \partial x_k} = \frac{\partial}{\partial x_j} \left(\sum_i r_i \frac{\partial r_i}{\partial x_k} \right) \\ &= \sum_i \underbrace{\frac{\partial r_i}{\partial x_j} \cdot \frac{\partial r_i}{\partial x_k}} + r_i \frac{\partial^2 r_i}{\partial x_j \partial x_k} \end{aligned}$$

$$\exists \sum_i \frac{\partial r_i}{\partial x_j} \cdot \frac{\partial r_i}{\partial x_k} = (y)_k (y)_j$$

$$\nabla^2 f = y^T y + \sum_i r_i (\nabla^2 r_i)$$

$\nabla^2 r_i$ often difficult and/or very expensive to compute. (sometimes impossible)

If $r \approx 0$ close to solution, or at least $\|r\| \ll \|y^T y\|$ then we neglect $\sum_i r_i (\nabla^2 r_i) \rightarrow$

Gauss-Newton method / model

Local model

In general J not square

Assume col.s of indep, otherwise $\tilde{n} < n$

Q_1, Q_1^T orthog.

proj. in $\langle \cdot, \cdot \rangle_2$

assoc. with $\|\cdot\|_2$

12.5

$$f(x_k + p) \approx m_k(p) = f_k + r^T J p + \frac{1}{2} p^T J^T J p$$

Newton step: $J^T J p = -J^T r$

Normal eq.s for $J p \approx -r \rightarrow \min_p \|J p + r\|_2$
 \rightarrow Linear Least Squares

$$J^{m \times n} \quad m > n$$

$$J : \mathbb{R}^n \rightarrow R(J) \subseteq \mathbb{R}^m$$

$$r \in \mathbb{R}^m = r_1 \mathbb{R}^n + \mathbb{R}^2 \left\{ \begin{array}{l} p_1 \in R(J) \\ p_2 \in R(J)^\perp \end{array} \right.$$

QR-decomp $J : J = Q \cdot R$
 \hookrightarrow orthog \hookrightarrow upper tri

$$R = \begin{bmatrix} \times & & & \\ & \times & & \\ & & \times & \\ & & & \times \\ & & & & 0 \end{bmatrix}$$

$$J = \begin{matrix} n & m-n \\ \begin{bmatrix} Q_1 & Q_2 \end{bmatrix} \end{matrix} \begin{bmatrix} R_1 \\ 0 \end{bmatrix} = Q_1 R_1$$

where $Q_1 Q_1^T = I_n \quad R(Q_1) = R(J)$

$$J p_1 = Q_1 Q_1^T r \quad (\text{part we can solve for})$$

$$\cancel{Q_1} (\mathbb{I} - Q_1 Q_1^T) r \quad (\text{residual})$$

$$Q_1^T (J p - r) = 0 \rightarrow \text{no comp. of}$$

residual in space $R(J)$ (otherwise better solution exists)

$$Q_1^T (Q_1 R_1 p - r) = 0 \Leftrightarrow R_1 p = Q_1^T r$$

Alternatively $y = U \Sigma V^T$

$$\begin{array}{ccc} [U_1 | U_2] & \begin{bmatrix} \sigma_1 & & \\ & \ddots & \\ & & \sigma_p & \\ & & & 0 \end{bmatrix} & \begin{bmatrix} V^T \\ \\ \\ \end{bmatrix} = \\ \begin{array}{c} m \times m \\ \\ \\ \end{array} & \begin{array}{c} m \times n \\ \\ \\ \end{array} & \begin{array}{c} n \times n \\ \\ \\ \end{array} \end{array}$$

$m \times n \quad n \times n \quad n \times n$

$$U_1 \Sigma_p V^T \quad (\text{allows for } \sigma_i \geq \sigma_p)$$

allows for some (or all) of σ_i to be zero.

Let $p \leq \min(m, n)$ s.t.

$$\sigma_1 \geq \sigma_2 \geq \dots \geq \sigma_p > \sigma_{p+1} = \dots = \sigma_n = 0$$

$$y = [u_1 \dots u_p] \begin{bmatrix} \sigma_1 & & \\ & \ddots & \\ & & \sigma_p \end{bmatrix} \begin{bmatrix} v_1^T \\ \vdots \\ v_p^T \end{bmatrix}$$

$$R(y) = \mathcal{S}\{u_1, \dots, u_p\}$$

$$p = V_p \Sigma_p^{-1} U_p^T r \quad \begin{array}{l} \text{pseudo} \\ p \text{ inverse} \end{array}$$

(least norm solution / gen. inverse)

Such solutions or approx. can be computed in a number of ways.

Damped GN: ~~iterative~~

p is (approx) sol. of LS. problem

$$x_{k+1} = x_k + \alpha_k p_k$$

$$U_p \Sigma_p V_p^T$$

also works if

$$m < n!$$

↙

where α_k typically found by backtracking alg.

→ First $\alpha_k = 1$ if Armijo cond. sat. accept, otherwise reduce α_k , until Armijo cond sat.

→ Book: if Wolfe conds satisfied and Y full rank (throughout R) and

$$\begin{aligned} \sigma_{\min}(Y) &\geq \gamma > 0 \\ \sigma_{\max}(Y) &\leq \beta < \infty \end{aligned}$$

then convergence guaranteed (std line search alg, Zoutendijk's condition and bound on $\cos \theta_k$)

If $r(x^*) = 0 \rightarrow$ quadr. conv.
otherwise \rightarrow (super) linear conv.

Even if $Y^T Y$ poss. singular, conv. can often be guaranteed if some, more elaborate, weaker cond. are satisfied.

$$p_{GN} = -V_p \Sigma_p^{-1} U_p^T r$$

$$-\nabla f^T p_{GN} = -r^T Y p$$

$$= +r^T U_p \Sigma_p \cancel{V_p^T V_p} \Sigma_p U_p^T r$$

$$= + (U_p^T r)^T \Sigma_p^2 (U_p^T r)$$

unif. bnd from 0 of $p \geq 1$

$$(U_p^T r)^T \Sigma_p^2 (U_p^T r) \geq \rho > 0 \text{ and } \sigma_p \geq \bar{\sigma} > 0 \\ \hookrightarrow \sigma_{\min} > 0$$

We can guarantee pos. def of approx.

Hessian : $(y^T y + dI)$ for any $d > 0$

This leads essentially to trust region approach or slight variant:

Levenberg-Marquardt method

$$p_{LM} = -(y^T y + dI)^{-1} y^T r$$

(Solve $(y^T y + dI)p = -y^T r$)

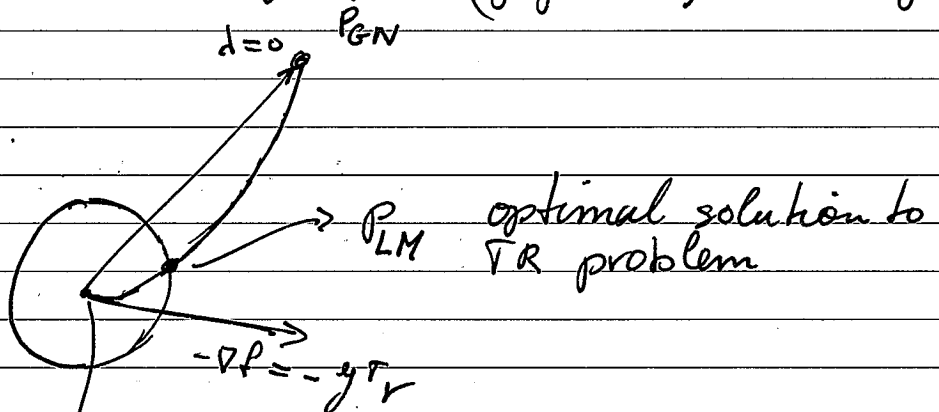
Choice of d important and similar to TR radius

$$\begin{cases} \min_p \frac{1}{2} \|y_k p + r_k\|_2^2 \\ \|p\|_2 \leq \Delta_k \end{cases}$$

~~Note~~ Assume $(y^T y)^{-1} y^T r = p_{GN}$ and

$$\|p_{GN}\| > \Delta$$

Consider $p(d) \leftarrow (y^T y + dI)^{-1} y^T r = -y^T r$



$$d \rightarrow \infty : (dI)p = -y^T r \Rightarrow p = -\frac{1}{d} y^T r$$

↳ LR gives solution balancing / weighting (neg) gradient and (Gauss) Newton step.

Compute in LS sense

$$p_{LS} = \arg \min_p \frac{1}{2} \left\| \begin{bmatrix} y \\ Vd \Sigma \end{bmatrix} p + \begin{bmatrix} r \\ 0 \end{bmatrix} \right\|_2^2$$

For implementation details, see pp. 259 and further

A convergence proof largely follows the std proof for TR methods.

~~Can~~ Comparison of Damped GN and LN.

$$y = \overset{n \times p}{U} \overset{p \times p}{\Sigma} \overset{p \times n}{V^T} \quad l \leq \min(m, n)$$

$$GN \rightarrow p_{GN} = -V_l \Sigma_l^{-1} U_l^T r \quad \sigma_l > \sigma_{l+1} = 0$$

$$(y^T y = V_l \Sigma_l^T \Sigma_l V_l^T, y^T r = V_l \Sigma_l^T U_l^T r)$$

$$\left(\begin{aligned} y^T y \cdot \frac{1}{\sigma_i^2} &= V_l \Sigma_l^{-2} V_l^T V_l \Sigma_l U_l^T r \\ y^T r &= V_l \Sigma_l^{-1} U_l^T r \end{aligned} \right)$$

Assume line search needed \rightarrow

$$p_k = \alpha_k p_{GN} = -\alpha_k V_l \Sigma_l^{-1} U_l^T r$$

If some σ_i (very) small, typically α_k very small

$$\text{Note } p_k = V_l \begin{pmatrix} \alpha_k & & \\ & \ddots & \\ & & \alpha_k \end{pmatrix} \Sigma_l^{-1} U_l^T r$$

$$= - \sum_{i=1}^l v_i \cdot \alpha_k \cdot \frac{u_i^T r}{\sigma_i}$$

α_k small, all comps damped equally strongly (indep of need) | LS damping applied component-wise

assuming we \leftarrow
compute least
norm solutions
of y singular

LM.

~~$$y = U_q \Sigma_q V_q^T$$~~

$$y = U_q \Sigma_q V_q^T$$

~~$$y^T y + dI = \dots$$~~

$$q = \min(m, n)$$

$$\Sigma_q = \text{diag}(\sigma_1, \dots, \sigma_q) \text{ and}$$

possibly some $\sigma_i = 0$ (assume $\sigma_1 \geq \sigma_2 \geq \dots$)

$$y^T y = V_q \Sigma_q^T \Sigma_q V_q^T$$

$$\hookrightarrow \begin{pmatrix} \sigma_1^2 & & \\ & \dots & \\ & & \sigma_q^2 \end{pmatrix} \text{ dim } q \times q$$

$$y^T y + dI = V_q (\Sigma_q^T \Sigma_q + dI) V_q^T$$

$$\begin{pmatrix} \sigma_1^2 + d & & \\ & \dots & \\ & & \sigma_q^2 + d \end{pmatrix} \begin{pmatrix} \sigma_1^2 & & \\ & \dots & \\ & & \sigma_q^2 \end{pmatrix}$$

$$(y^T y + dI)^{-1} y^T r =$$

$$= V_q (\Sigma_q^T \Sigma_q + dI)^{-1} V_q^T V_q \Sigma_q^T U_q^T r$$

$$= - \sum_{i=1}^q v_i \frac{\sigma_i}{\sigma_i^2 + d} u_i^T r$$

$$= - \sum_{i=1}^q v_i \frac{\sigma_i^2}{\sigma_i^2 + d} \cdot \frac{u_i^T r}{\sigma_i}$$

 \hookrightarrow damping coeff.s

damping of indiv. components depends on magn. of d vs σ_i

 $\sigma_i^2 \gg d \rightarrow$ very little damping

 $\sigma_i^2 = d \rightarrow \frac{1}{2}$
 $\sigma_i^2 \ll d \rightarrow$ very strong damping

$$\min_{x \in \mathbb{R}^n} f(x) \quad \text{subj to} \quad \begin{cases} c_i(x) = 0 & i \in E \\ c_i(x) \geq 0 & i \in I \end{cases}$$

f, c_i all smooth, real valued on subset of \mathbb{R}^n

I, E finite index sets

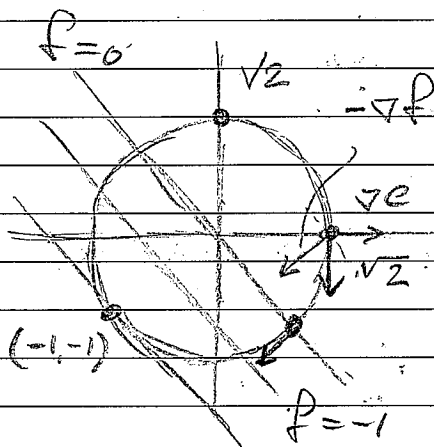
$$\Omega = \{x \in \mathbb{R}^n \mid c_i(x) \geq 0, i \in I, c_i(x) = 0, i \in E\}$$

$\hookrightarrow \Omega$ bounded by piecewise smooth curves/surfaces

Active set $A(x)$ at any feasible x consists of set of indices of constr. satisfied as equality constraints.

$$A(x) = \sum \cup \{i \in I \mid c_i(x) = 0\}$$

Ex. 12.1 $\min_{x_1, x_2} x_1 + x_2 \quad \text{s.t.} \quad x_1^2 + x_2^2 - 2 = 0$



$$x_1 + x_2 = c \quad \rightarrow \quad \begin{aligned} x_1 &= c - x_2 \\ x_2 &= c - x_1 \end{aligned}$$

x : stay on constr. curve

$$\begin{aligned} c(x) &= 0 \quad \text{and} \\ c(x+s) &= c(x) + \nabla c(x)^T s + \mathcal{O}(\|s\|^2) \\ &= 0 \end{aligned}$$

for small s : $\nabla c^T s = 0$

To reduce f : $\nabla f^T s < 0$

Unless ∇f and ∇c parallel, we can find direction s which reduces f along constraint curve/surface.

Altern. if such λ does not exist, we satisfy cond. s for (constr.) local extremum
 \rightarrow Lagrange multiplier method.

$$L = f(x) + \lambda c(x) \quad \rightarrow$$

$$\text{stat. point } L : \begin{cases} \frac{\partial L}{\partial x} = \nabla_x f + \lambda \nabla_x c = 0 \\ \frac{\partial L}{\partial \lambda} = c(x) = 0 \end{cases}$$

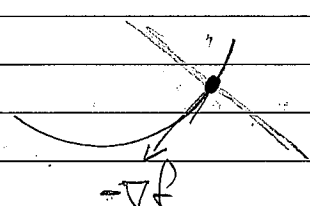
First cond: $\nabla f = -\lambda \nabla c$ parallel

$$\begin{cases} 1 + \lambda(2x_1) = 0 & 2\lambda x_1 = -1 \Rightarrow \lambda x_1 = -\frac{1}{2} \\ 1 + \lambda(2x_2) = 0 & \lambda x_2 = -\frac{1}{2} \\ x_1^2 + x_2^2 = 2 & \end{cases}$$

$$\frac{1}{4} x_1^2 + \frac{1}{4} x_2^2 = \frac{1}{4} 2d^2 \rightarrow d^2 = \frac{1}{4} \rightarrow d = \pm \frac{1}{2}$$

$$\begin{aligned} d = \frac{1}{2} &\rightarrow x_1 = x_2 = -1 \\ d = -\frac{1}{2} &\rightarrow x_1 = x_2 = 1 \end{aligned} \quad \left. \begin{array}{l} \text{stat. points} \\ \rightarrow \text{max} \\ \rightarrow \text{min} \end{array} \right\}$$

$$\min x_1 + x_2 \quad \text{s.t.} \quad x_1^2 + x_2^2 \leq 2 \rightarrow 2 - x_1^2 - x_2^2 \geq 0$$



Still need $\nabla f^T s < 0$

but only $c_1 + \lambda \nabla c^T s \geq 0$
 (to first order)

∇f x in interior of region $c(x) \geq 0$

any small enough step satisf. 2 crit.

So, only need $\nabla f^T s < 0$

$$\nabla f \nabla P \neq 0 \rightarrow s = -\alpha \nabla P$$

If x on bnd situation as before.

$$R = f + dc$$

Opt. (first order \rightarrow stat. pt) :

$$\nabla_x R = \nabla f + d \nabla c = 0 \text{ for some } d \geq 0$$

and $dc(x) = 0$ (either $d=0$ or $c=0$)
 \hookrightarrow complementarity condition

$\begin{cases} c=0 & d \text{ can be anything, constr. active} \\ & \text{and need } \nabla c \text{ parallel to } \nabla f \end{cases}$

$c > 0 \quad d = 0$ so, that constr. has no effect on (local) solution

Need ^{robust} ~~good~~ methods to come up with useful search directions taking both obj. function and constraints into account.

$\{z_k\}$ feasible sequence approaching x if $z_k \in \Omega$ for all k suff. large and $z_k \rightarrow x$.

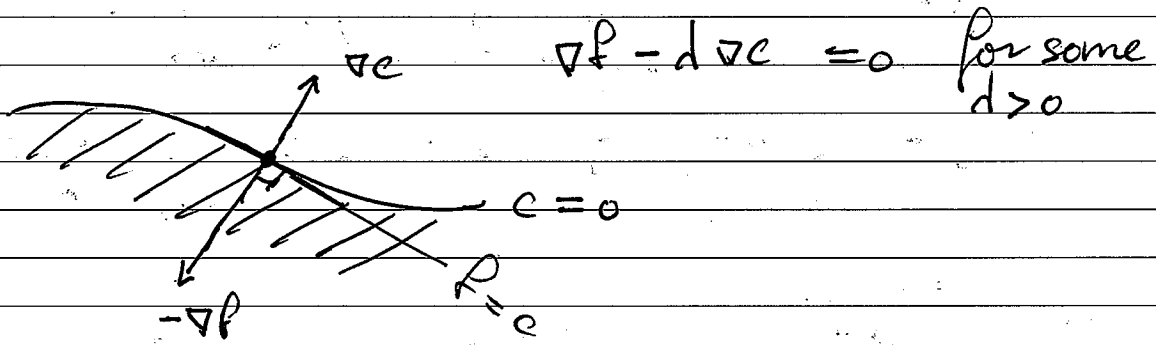
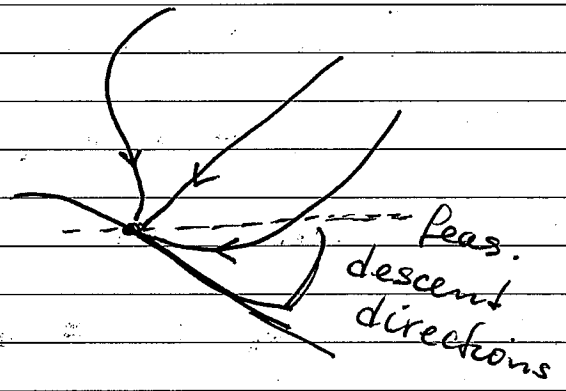
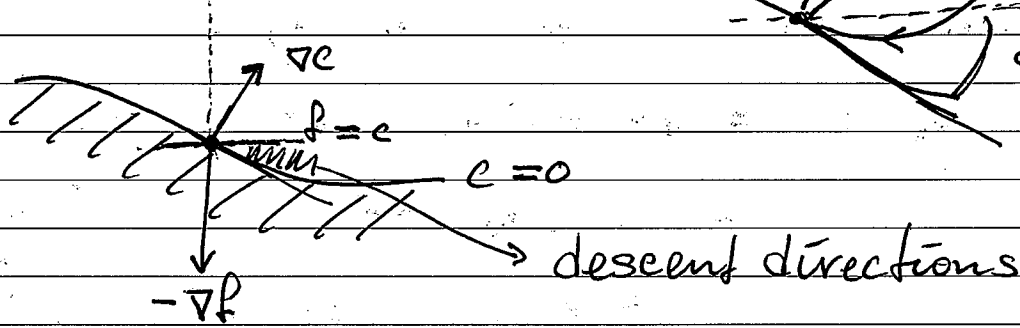
We can define local solution as point x^* s.t.

all feasible sequences approaching x^* have $f(z_k) \geq f(x^*)$ for all suff. large k .

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$$\min f(x) \quad \text{subj. } c_i \geq 0$$

$$\begin{cases} \nabla f - d \nabla c = 0 & d \geq 0 \\ c \geq 0 \end{cases}$$



Def 12.2
 Ω feasible region

For a set F is a cone if
 $x \in F \Rightarrow \alpha x \in F$
 $\forall \alpha > 0$

Def 12.3

The vector d is called a tangent (vector) to Ω at x if there are a feasible seq. $\{z_k\}$ approaching x and a seq. of pos. scalars $\{t_k\}$ w. $t_k \rightarrow 0$ s.t.

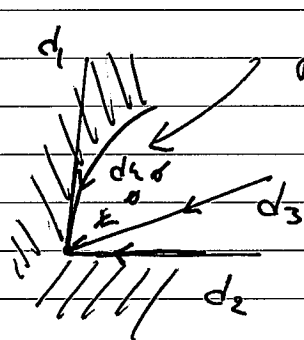
$$\lim_{k \rightarrow \infty} \frac{z_k - x}{t_k} = d$$

The set of all tangents to Ω at x is called the tangent cone and is denoted by $T_{\Omega}(x)$

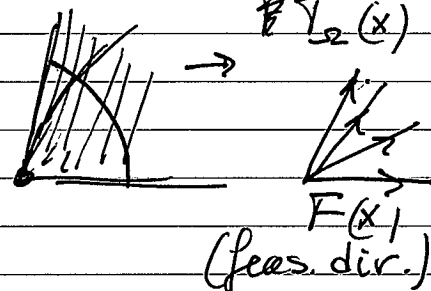
Given feas. point x and active constr. set $A(x)$, the set of linearized feas. directions $F(x)$ is

$$F(x) = \left\{ d \mid \begin{array}{l} d^T \nabla c_i = 0 \quad \forall i \in \mathcal{E} \\ d^T \nabla c_i \geq 0 \quad \forall i \in \mathcal{I} \cap A(x) \end{array} \right\} \text{ and } \downarrow \text{ active ineq. constr.}$$

(why compl. def? \rightarrow tangent cone depends on geometry of Ω , not a choice of alg. specification.)



feasible domain



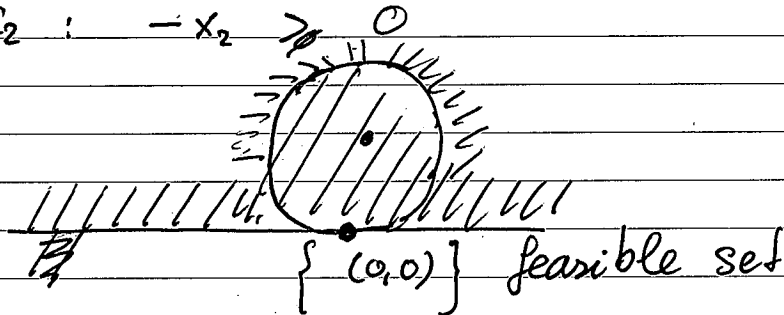
$$(x_1-1)^2 + (x_2-1)^2 \leq 1 \leftarrow$$

$$x_2 \leq 0 \leftarrow$$

$$c_1 : 1 - (x_1-1)^2 - (x_2-1)^2 \geq 0$$

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$$c_2 : -x_2 \geq 0$$



$$\begin{cases} d^T \nabla c_1 = 0 \\ d^T \nabla c_2 = 0 \end{cases}$$

at $(0,0)$ true for $d = \begin{pmatrix} x \\ 0 \end{pmatrix}$

but there is no direction to move

problem is that constr. are dependent.

Def 12.4

Given pt x and active set $A(x)$ the

linear indep. constraint qualification

(LICQ) holds if the set of active

constraint grads $\{\nabla c_i \text{ for } i \in A(x)\}$

is linearly indep.

~~When~~

Relative to linearized constraints,
what are necessary/suff. conditions
of optimality?

Consider Lagrangian

$$L(x, d) = f(x) - \sum_{i \in \mathcal{I} \cup \mathcal{E}} d_i c_i(x)$$

First-order necessary conditions \rightarrow

(nec. cond.s approx. to first order)

Theo ~~Def~~ 12.1

Suppose that x^* is a local solution of constr. min. problem in std form, that f and (all) c_i are cont. diff. and that LICQ holds at x^* . Then there is a Lagr. multiplier vector λ^* such that at (x^*, λ^*)

$$\begin{aligned} \nabla h(x^*, \lambda^*) &= 0 \\ c_i(x^*) &= 0 & \forall i \in \mathcal{E} \\ c_i(x^*) &\geq 0 & \forall i \in \mathcal{I} \\ \lambda_i^* &\geq 0 & \forall i \in \mathcal{I} \\ *1 \quad \lambda_i^* c_i(x^*) &= 0 & \forall i \in \mathcal{E} \cup \mathcal{I} \end{aligned}$$

Karush-Kuhn-Tucker conditions (KKT)

*1 complementarity conditions:

if $c_i(x^*) > 0$ then $d_i = 0$ (no effect) ^{constr.}

if $c_i(x^*) = 0$ then $d_i = 0$ or nonzer

\hookrightarrow pref. $d_i \neq 0$ otherwise some type of indeterminacy \rightarrow active constr. does not influence Lagr. (locally)

Def 12.5

Strict Complementarity

Given local sol. x^* and λ^* satisfying 1st order nec KKT cond.s, strict complementarity cond. holds if exactly one of $d_i^* = 0$ or $c_i(x^*) > 0$ holds for

each $i \in \mathcal{I}$. In other words, $d_i^* > 0$ for each $i \in \mathcal{I} \cap \mathcal{A}(x^*)$

When LICQ holds the (optimal) d^* at local sol. x^* is unique

$$\text{Def } A(x^*) = \left[\nabla C_i(x^*)^T \right]_{i \in A(x^*)}$$

Lemma: let x^* be feasible point. Then

$$(1) \quad T_{\Omega}(x) \subset F(x)$$

$$(2) \quad \text{If LICQ satisfied at } x, \text{ then } F(x) = T_{\Omega}(x)$$

Proof: see book

Theo 12.3

(necessary cond)

If x^* is local solution, then

$$\nabla f(x^*)^T d \geq 0 \text{ for all } d \in T_{\Omega}(x^*)$$

Proof:

Assume not true: $\exists d, \nabla f(x^*)^T d < 0$

$$z_k = x^* + t_k d + o(t_k)$$

Consider $\{z_k\}, \{t_k\}$ s.t. $\lim_{k \rightarrow \infty} \frac{z_k - x^*}{t_k} = d$
 \rightarrow feas. seq.

$$\begin{aligned} f(z_k) &= f(x^*) + (z_k - x^*)^T \nabla f(x^*) + o(\|z_k - x^*\|) \\ &= f(x^*) + t_k d^T \nabla f(x^*) + o(t_k) \end{aligned}$$

Since $d^T \nabla f(x^*) < 0$ (constant)

and $t_k > 0$
for k large enough $f(z_k) < f(x^*)$

So, x^* not local solution (min in feas. set)

Lemma 12.4 (Farkas)

Let $B \in \mathbb{R}^{n \times m}$ and $C \in \mathbb{R}^{n \times p}$ and define cone

$K = \{ By + Cw \mid y \geq 0 \}$. Then ~~either~~ exactly
for any $g \in \mathbb{R}^n$
one of ~~the~~ two cases holds

(either) (1) $g \in K$

(or) (2) $\exists d \in \mathbb{R}^n : g^T d < 0, B^T d \geq 0, C^T d = 0$

comp. wise

~~A~~

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Assume f and e_i twice cont. diff.

Assume (x^*, d^*) sat. KKT cond.s

Def. critical cone $C(x^*, d^*) =$

$$\{w \in F(x^*) \mid \nabla c_i(x^*)^T w = 0 \text{ for all } i \in A(x^*) \cap \Sigma \text{ with } d_i^* > 0\}$$

$$w \in C(x^*, d^*) \Leftrightarrow \begin{cases} \nabla c_i(x^*)^T w = 0 & \forall i \in \Sigma \\ \nabla c_i(x^*)^T w = 0 & \forall i \in A(x^*) \cap \Sigma, w_i d_i^* > 0 \\ \nabla c_i(x^*)^T w \geq 0 & \forall i \in A(x^*) \cap \Sigma, w_i d_i^* = 0 \end{cases}$$

$$\rightarrow d_i^* w_i^T \nabla c_i(x^*) = 0 \quad \forall i \in \Sigma \cup \Gamma$$

$$w \in C \Rightarrow w^T \nabla f(x^*) = \sum_{i \in \Sigma \cup \Gamma} d_i^* w_i^T \nabla c_i(x^*) = 0$$

\hookrightarrow

direction for which first derivative information does not give us info. whether f increases or decreases

Theo 12.5 Suppose x^* is local sol. and the LICQ

is satisfied. Let d^* be the Lagr. mult for which

KKT satisfied. Then

$$w^T \nabla_{xx} L(x^*, d^*) w \geq 0 \text{ for all } w \in C(x^*, d^*)$$

(second or necessary cond.)

Second order suff. cond.s

x^* feasible, KKT cond. hold

$$w^T \nabla_{xx} L(x^*, d^*) w > 0 \text{ for all } w \in C, w \neq 0$$

x^* is strict local minimizer